Declarative program development in Prolog with GUPU

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- Programming environment for beginners
- New program development **process** specification & implementation phase
- All phases are supported by diagnostic facilities
- Emphasizing notion of relation

GUPU

GUPU

Gesprächs

 \mathbf{G} esprächs \mathbf{u} nterstützende

 $\mathbf{G} \mathbf{e} \mathbf{s} \mathbf{p} \mathbf{r} \mathbf{\ddot{a}} \mathbf{c} \mathbf{h} \mathbf{s} \mathbf{u} \mathbf{n} \mathbf{t} \mathbf{e} \mathbf{r} \mathbf{s} \mathbf{\ddot{u}} \mathbf{t} \mathbf{z} \mathbf{e} \mathbf{n} \mathbf{d} \mathbf{e}$

\mathbf{P} rogrammierübungs

1, 2, 3, 4,

 \mathbf{G} esprächs \mathbf{u} nterstützende

 \mathbf{P} rogrammierübungs \mathbf{u} mgebung

1, 2, 3, 4, 5,

Gesprächsunterstützende

Programmierübungs**u**mgebung environnment

1, 2, 3, 4, 5, 6,

 \mathbf{G} esprächsunterstützende

 \mathbf{P} rogrammierübungs \mathbf{u} mgebung cours $\stackrel{de}{\leftarrow}$ environnment

1, 2, 3, 4, 5, 6, 7,

Gesprächsunterstützende

1, 2, 3, 4, 5, 6, 7, 8,

1, 2, 3, 4, 5, 6, 7, 8, 9,

GUPU - explication

1, 2, 3, 4, 5, 6, 7, 8, 9, 10,

1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11,

GesprächsunterstützendeProgrammierübungsumgebungconversations $\stackrel{\text{des}}{\leftarrow}$ supportant \leftarrow programmation $\stackrel{\text{de}}{\leftarrow}$ cours $\stackrel{\text{de}}{\leftarrow}$ environnmentConversation \rightarrow supporting

1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12,

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- Used since 1992
- \bullet Under continual development since 1991
- Original motivation: realize courses with a large number of students
- Eases assessment (marking) instantaneous, automated pre-marking
- General attitude: Mark now, don't delay it unto the end
- \bullet 9 weeks/about 80 (small) exercises
- Flexible low cost system for deadlines
- \bullet Simple to use very simple interaction mode
- Consistent view of program
- Useless notions absent (files, shells, overlapping windows etc.)
- Side effect free. Pure, monotone subset of Prolog including constraints
- currently trilingual (German, French, English)

 $1,\,2,\,3,\,4,\,5,\,6,\,7,\,8,\,9,\,10,\,11,\,12,\,13,\,14,\,15,\,16$

- easy to confuse with procedures
- Algorithm = Logic + Control.

• Conclusion:

3

1,

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What is the input/output of a predicate?

Answer:

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• Naming: avoid imperative names — helps somewhat, but soon forgotten.

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- easy to confuse with procedures
- Algorithm = Logic + Control. Often: Algorithm \approx Control \Rightarrow Logic ≈ 0
- input/output: What is the input/output of a predicate?

Answer: There is no input/output — not helpful.

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How to focus on the declarative properties?

 $1,\,2,\,3,\,4,\,5,\,6,\,7,\,8,\,9,\,10,\,11,\,12$

Extreme Programming

Lightweight, agile method developed by Kent Beck for Smalltalk. Practice to Code Unit Test First *Test program into existence!*

- All code must have unit tests.
- All code must pass all unit tests before it can be released.
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because of **logical variables**.

- less specific: \leftarrow alldifferent(Xs).
- higher coverage: $\not\leftarrow$ all different([X,X|_]).

... than in traditional languages

- Tests: assertions
 - Positive assertions: \leftarrow Goal should succeed
 - Negative assertions: $\not\leftarrow$ Goal should fail
- Close integration: Tests are written *into the program text*
- All predicates must have assertions
- Errors are signaled immediately within the program text, explanations based on *slicing* are offered
- Adding further assertions very easy
 - Duplicate and modify existing assertion
 - Offered by diagnostic facilities
- Tests are run very often: Upon every saving, all assertions are tested

1. Start with the least specific test.

 $\leftarrow all different (Xs). There is at least a single solution, Xs is anything$

- 2. Estimate cardinality of *minimal* possible set of answer substitutions. If infinite, goal *must not* terminate. *♀* alldifferent(Xs), false.
- 3. Go further to more specific tests. $\leftarrow Xs = [_,_], all different(Xs).$
- 4. For every positive assertion, find a similar negative assertion. $\not\leftarrow Xs = [1,1]$, all different (Xs).
- 5. Generalize negative assertions as much as possible. $\not\leftarrow Xs = [X,X]$, all different (Xs).
- 6. Specialize positive assertions as much as possible. $\leftarrow Xs = [1,2]$, all different (Xs).

But, one problem remains...

Biggest obstacles to testing prior to coding:

- Cumbersome to write tests containing lots of data
- Incorrect tests slows development
- No motivation to write tests since they might be wrong
- Adjusting tests to the program

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Conclusion:

• Attention span too large for beginners

Solution:

• Put learner into the position of testing predicates prior to writing them

Reference implementation — testing the tests

Assertions are tested against reference implementation. Reference implementation is considered correct for

- unconditional success (no pending constraints)
- finite failure

Reference implementation is ignored for:

- \bullet (implementation related reference implementation not perfect)
 - exceptions
 - $-\operatorname{computation}$ takes too long/loops
 - $-\operatorname{conditional}$ success with constraints that cannot be resolved
- (specification related relation is under-specified) Signaled as exceptions or constraints. E.g.: child_of/2

\Rightarrow All procedural issues are ignored.

Marking system already counts correct assertions.

- reason: there is a solution
- \bullet show solution in the form of a positive assertion
- try to make assertion as specific as possible
 - show binding (answer substitution)
 - $-\operatorname{try}$ to ground remaining variables with constants any 1, \ldots
 - $\not\leftarrow \mathrm{Xs} = [_,_,_|_], \text{ all different}(\mathrm{Xs}).$
 - @@ % != Should be positive!
 - @@% Even this specialized assertion should be true
 - $@@ \leftarrow Xs = [any0,any1,any2], all different(Xs).$
 - $-\operatorname{try}$ to ground fd-variables with some values

- reason: there is no solution
- \bullet show a generalized goal in form of a negative assertion
- \bullet try to generalize assertion to better localize the error
 - $\leftarrow all different([a, b, c, d, c, f]).$
 - @@ % != Should be negative!
 - @@%@ Generalized negative assertion
 - $\textcircled{0} \textcircled{0} \not\leftarrow all different([_,_,c,_,c,_]).$
 - @@ % @ Further generalization
 - $\textcircled{0} \textcircled{0} \not\leftarrow \text{alldifferent}([_,_,V0,_,V0,_]).$
 - @@ % @ Generalization by goal replacement
 - $@@ \not\leftarrow all different([V0,_-,V0|_-]).$
 - $\leftarrow [a,b,c,d] = [a,b,e,d].$
 - @@ % != Fails as it should!
 - @@%@ Generalized negative assertion
 - $\textcircled{0} \textcircled{0} \not\leftarrow [_,_,c|_] = [_,_,e|_].$
 - @@%@ Generalization using dif/2

 $code_inconnu/2$:

- Nothing is said about the relation except that you will only get information about it via assertions
- Relation defined differently for everyone

Effects of testing with reference implementation

- + test coverage significantly better
- + more than twice as many assertions are written
- + almost no incorrect programs (i.e. automatic marking almost perfect)
- + students consider (and question) the example statements more closely
- + almost no student questions concerning example statements (most frequent question previously: *What is the output?*)
- + (the very few) questions focus rather on the specification itself
- + more fun due to fast response

After coding: reading of programs

- traditional readings: declarative and procedural
- selective readings: use transformations to obtain slices (fragments)
 - generalization: delete goals

father(Father) \leftarrow * $\overline{\text{male}(\text{Father})},$ child_of(_C, Father).

specialization: add goals (false/0: failure slice).

 $\begin{array}{l} \hline married_to(Husband, Spouse) \leftarrow false, \\ \hline husband_spouse(Husband, Spouse). \\ married_to(Spouse, Husband) \leftarrow \\ husband_spouse(Husband, Spouse). \\ \end{array}$

- + eases reading of larger programs
- + remains close to source code, simple presentation by hiding parts
- + no new formalism like proof trees, traces
- + works also with incomplete constraints

insufficiency (unexpected success): maximal failing generalization explains *data inconsistency* and *modeling errors*

- **incorrectness (unexpected success):** maximal specialization (with false/0) that succeeds
- **non-termination:** maximal non-terminating specialization
- Common properties:
- + error in fragment implies error in original program
- + visible fragment has to be changed
- + no user-interaction (\Rightarrow no debugging errors possible!)
- ? slicing or program modification ?

- side effect free visualization of answer substitutions
- general form: \leftarrow Viewer \ll Goal.
- $\bullet \ll$ can only be used within assertions, not allowed in rules
- most viewers are implemented side effect free within GUPU
- very few elementary viewers text, postscript

Problems searching for explanations of unexpected failure

- non-termination because of generalized fragments \rightarrow analyse termination (cTI)
- complexity: sub-problem already NP-hard, no approximation possible (Monotone Minimum Satisfying Assignment, Umans 1999)
 → search local minima, one by one (one test per line)
- labeling for generalized fragments often very expensive \rightarrow adopt labeling strategy

Similar sub-problem: Explanations in PPC (Narendra Jussien)

- generalization of (dynamic) constraint system
- much more constraints (at runtime) than (static) program points
 - more costly
 - less readable but contains more information
- uses a search interlaced with labeling (very interesting!)