

See-code und Code-Erzeugung in Gforth 1.0

M. Anton Ertl, TU Wien

See und simple-see

: bar if 5 then ;

see bar

: bar

IF #5

THEN ;

simple-see bar

\$3F9A20A540 ?branch

\$3F9A20A548 <bar+\$20>

\$3F9A20A550 lit

\$3F9A20A558 #5

\$3F9A20A560 ;s

See-code

see-code bar

\$3F9A20A540 ?branch 1->1

\$3F9A20A548 <bar+\$20>

3F99EBAC54: addi sp,sp,8
3F99EBAC56: mv spB,spA
3F99EBAC58: ld spA,\$0(sp)
3F99EBAC5C: ld a5,\$0(ip)
3F99EBAC60: addi a6,ip,\$8
3F99EBAC64: bnez spB,3F99EBAC6E
3F99EBAC66: ld a4,\$0(a5)
3F99EBAC68: addi ip,a5,\$8
3F99EBAC6C: jr a4
3F99EBAC6E: addi ip,a6,\$8

\$3F9A20A550 lit 1->1

\$3F9A20A558 #5

3F99EBAC72: sd spA,\$0(sp)
3F99EBAC76: addi sp,sp,-8
3F99EBAC78: ld spA,\$0(ip)
3F99EBAC7C: addi ip,ip,10
\$3F9A20A560 ;s 1->1
3F99EBAC7E: ld a6,\$0(rp)
3F99EBAC82: addi rp,rp,8
3F99EBAC84: addi ip,a6,\$8
3F99EBAC88: ld a4,\$-8(ip)
3F99EBAC8C: jr a4

Non-relocatable primitives

```
: erase 0 fill ;          see-code erase
                          $3F98F9E540 lit      1->1
                          $3F98F9E548 #0
                          3F98C4EC54:      sd      spA,$0(sp)
                          3F98C4EC58:      addi    sp,sp,-8
                          3F98C4EC5A:      ld      spA,$0(ip)
                          3F98C4EC5E:      addi    ip,ip,10
                          3F98C4EC60:      ld      a4,$-8(ip)
                          3F98C4EC64:      jr      a4
                          $3F98F9E550 fill
                          $3F98F9E558 ;s     1->1
                          3F98C4EC66:      ld      a6,$0(rp)
                          3F98C4EC6A:      addi    rp,rp,8
                          3F98C4EC6C:      addi    ip,a6,$8
                          3F98C4EC70:      ld      a4,$-8(ip)
                          3F98C4EC74:      jr      a4
```

Static Superinstructions

```
: foo < if 5 then ;
```

\$3F98F9E588 < ?branch	1->1	\$3F98F9E5A0 lit	1->1
\$3F98F9E590 ?branch		\$3F98F9E5A8 #5	
\$3F98F9E598 <foo+\$28>		3F98C4EC9C:	sd spA,\$0(sp)
3F98C4EC76:	ld a4,\$8(sp)	3F98C4ECA0:	addi sp,sp,-8
3F98C4EC7A:	ld a5,\$8(ip)	3F98C4ECA2:	ld spA,\$0(ip)
3F98C4EC7E:	addi sp,sp,10	3F98C4ECA6:	addi ip,ip,10
3F98C4EC80:	addi a6,ip,\$10	\$3F98F9E5B0 ;s	1->1
3F98C4EC84:	blt a4,spA,\$3F98C4EC94	3F98C4ECA8:	ld a6,\$0(rp)
3F98C4EC88:	ld a4,\$0(a5)	3F98C4ECAC:	addi rp,rp,8
3F98C4EC8A:	ld spA,\$0(sp)	3F98C4ECAE:	addi ip,a6,\$8
3F98C4EC8E:	addi ip,a5,\$8	3F98C4ECB2:	ld a4,\$-8(ip)
3F98C4EC92:	jr a4	3F98C4ECB6:	jr a4
3F98C4EC94:	ld spA,\$0(sp)		
3F98C4EC98:	addi ip,a6,\$8		

Constant-based optimization

```
: foo 2 pick ;
```

```
see foo
```

```
: foo
```

```
  third ;
```

```
see-code foo
```

```
$3F7ADE05F8 third 1->1
```

```
3F7AA90CAE: ld      a5,$10(sp)
```

```
3F7AA90CB2: addi   ip,ip,8
```

```
3F7AA90CB4: addi   sp,sp,-8
```

```
3F7AA90CB6: sd     spA,$8(sp)
```

```
3F7AA90CBA: mv     spA,a5
```

```
$3F7ADE0600 ;s 1->1
```

```
3F7AA90CBC: ld     a6,$0(rp)
```

```
3F7AA90CC0: addi   rp,rp,8
```

```
3F7AA90CC2: addi   ip,a6,$8
```

```
3F7AA90CC6: ld     a4,$-8(ip)
```

```
3F7AA90CCA: jr     a4
```

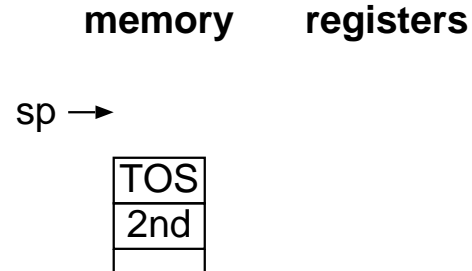
Stack Caching

```

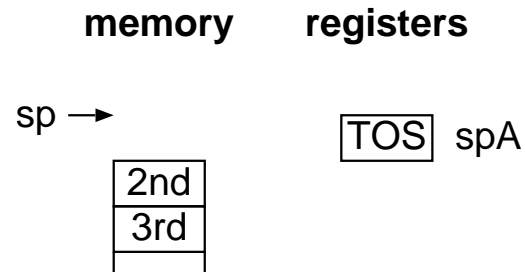
: foo swap ! ;
see-code foo
$3F7ADE0540 swap      1->2
    3F7AA90C54:  ld      spB,$8(sp)
    3F7AA90C58:  addi    ip,ip,8
    3F7AA90C5A:  addi    sp,sp,8
$3F7ADE0548 !        2->0
    3F7AA90C5C:  sd      spA,$0(spB)
    3F7AA90C60:  addi    ip,ip,8
$3F7ADE0550 ;s      0->1
    3F7AA90C62:  ld      spA,$8(sp)
    3F7AA90C66:  addi    sp,sp,8
    3F7AA90C68:  ld      a6,$0(rp)
    3F7AA90C6C:  addi    rp,rp,8
    3F7AA90C6E:  addi    ip,a6,$8
    3F7AA90C72:  ld      a4,$-8(ip)
    3F7AA90C76:  jr      a4

```

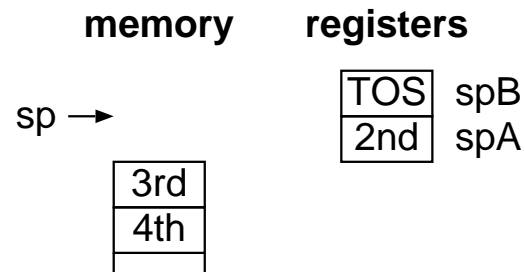
State 0



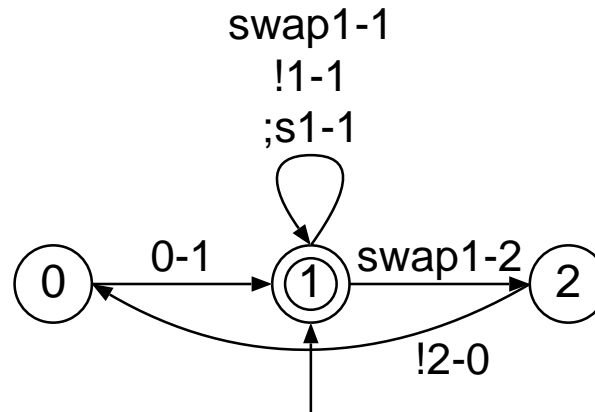
State 1



State 2



Mögliche Implementation



- Idealerweise eine Sequenz auf einmal
Aber: Decompiler schaut sich lieber jedes primitive einzeln an
- Möglicherweise mehrdeutige Sequenzen
Und wenn schon!
Bei aktueller Codeerzeugung vermutlich eindeutig
- Kostet keinen Extra-Speicher
Aber dafür Implementationsaufwand

Tatsächliche Implementierung

- Compiler speichert für jedes compilierte Primitive:
Startadresse und Länge des native code
Welches primitive oder superinstruction ist es
Anfangs- und Endzustand des Stack Cache
- Der Decompiler braucht nur mehr nachschauen
- Kostet 16 Bytes pro compiliertem primitive
- ca. 500KB für das Gforth-Image
fast soviel wie der Platz für den native code
- Einfacher zu implementieren

Zusammenfassung

- see-code zeigt native code
- zeigt static superinstructions
- zeigt stack caching
- Implementierung mit Zusatzinformation durch den Compiler