

## Limited Address Range Architecture for Reducing Code Size in Embedded Processors

Qin Zhao, Bart Mesman, Henk Corporaal  
Eindhoven University of Technology, The Netherlands  
Philips Research Laboratories, The Netherlands

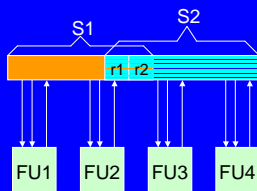
- Introduction
- The proposed architecture: LAR
- Sequential code generation for LAR
  - Annotated Conflict Graph(ACG)
- The integrated approach
  - Annotated Worst-Case Conflict Graph(AWCCG)
- Experimental results
- Conclusions and future work

### • Introduction

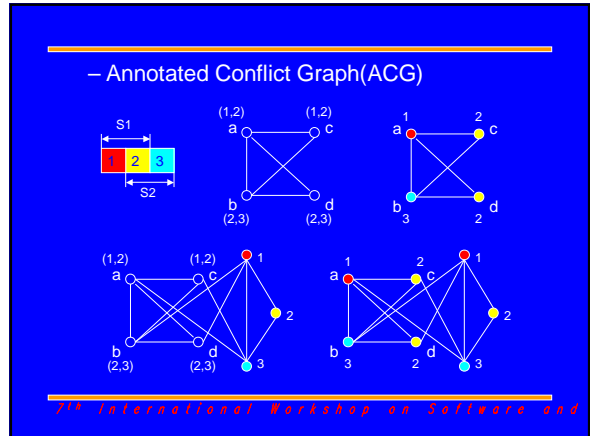
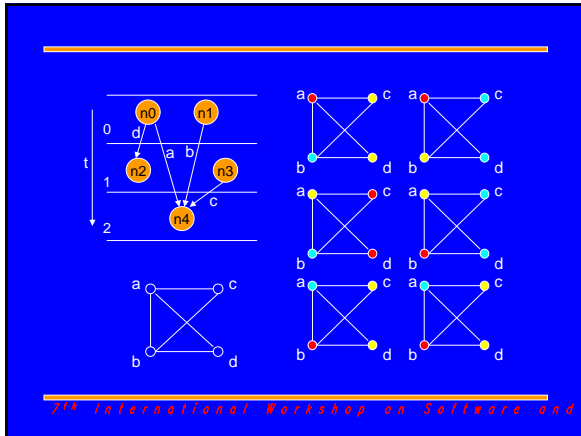
- Code size, power consumption of embedded cores must be small since they are on chip
- Irregularities in architectures
  - Difficult for efficient code generation
- Clustered register file vs. central register file
  - Advantage: small code size, power consumption
  - Disadvantage: extra hardware, copy operations
- Phase coupling in code generation
  - Sequential phases may generate inefficient code
  - Integrated approach potentially offers better solutions

### • Introduction

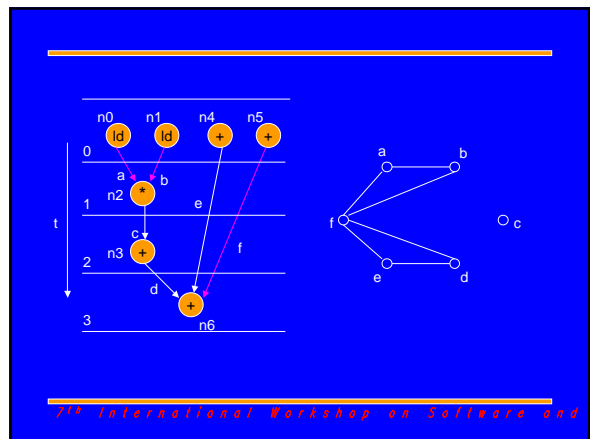
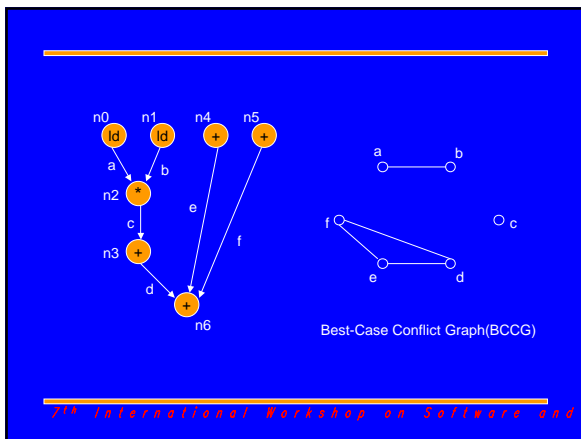
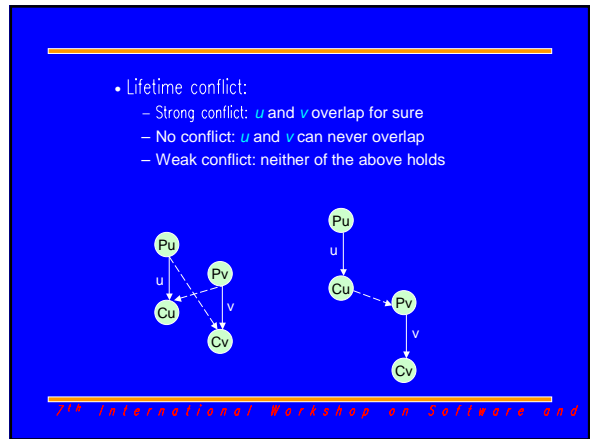
- The proposed architecture: LAR
- Sequential code generation for LAR
  - Annotated Conflict Graph(ACG)
- The integrated approach
  - Annotated Worst-Case Conflict Graph(AWCCG)
- Experimental results
- Conclusions and future work

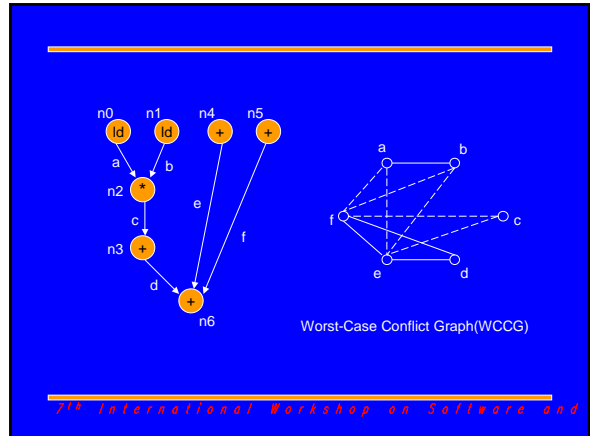
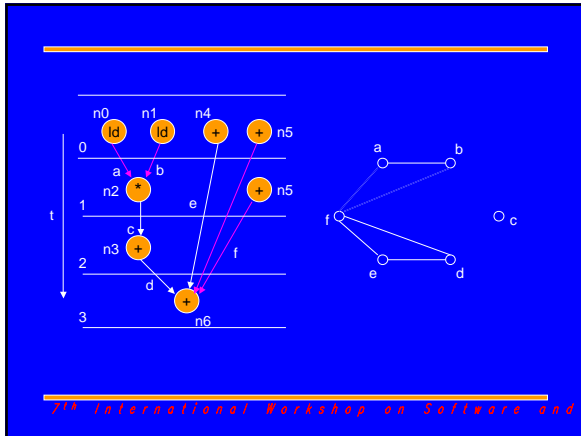


- Introduction
- The proposed architecture: LAR
- Sequential code generation for LAR
  - Annotated Conflict Graph(ACG)
- The integrated approach
  - Annotated Worst-Case Conflict Graph(AWCCG)
- Experimental results
- Conclusions and future work



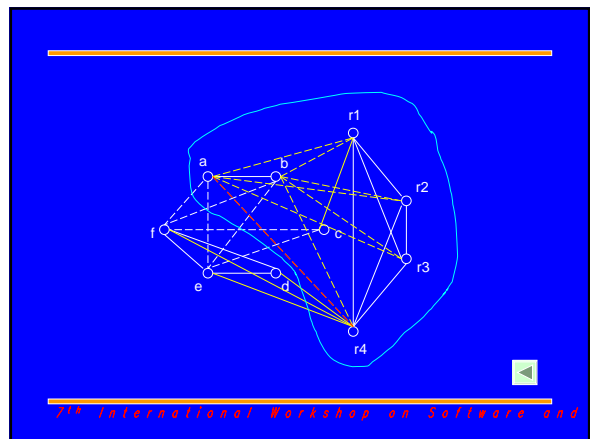
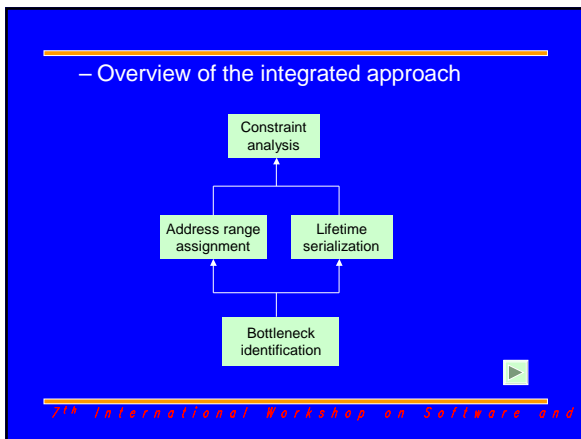
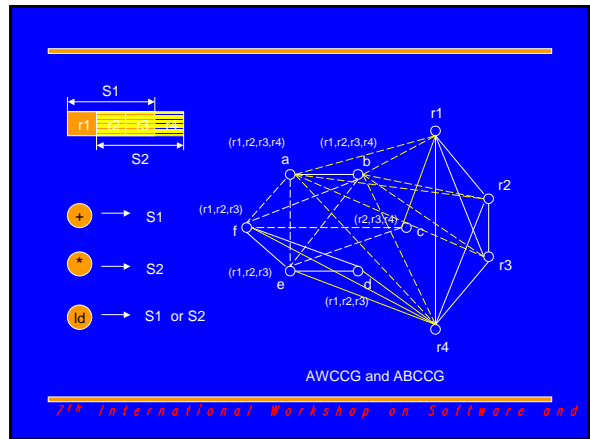
- Introduction
- The proposed architecture: LAR
- Sequential code generation for LAR
  - Annotated Conflict Graph(ACG)
- The integrated approach
  - Annotated Worst-Case Conflict Graph(AWCCG)
- Experimental results
- Conclusions and future work

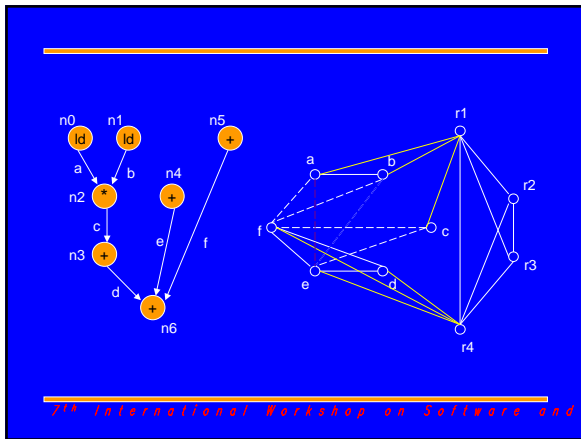
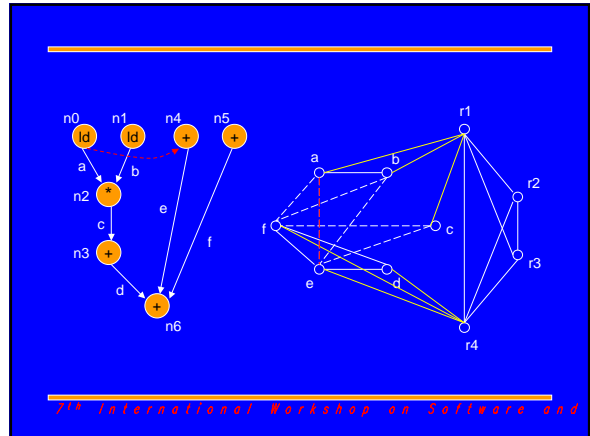
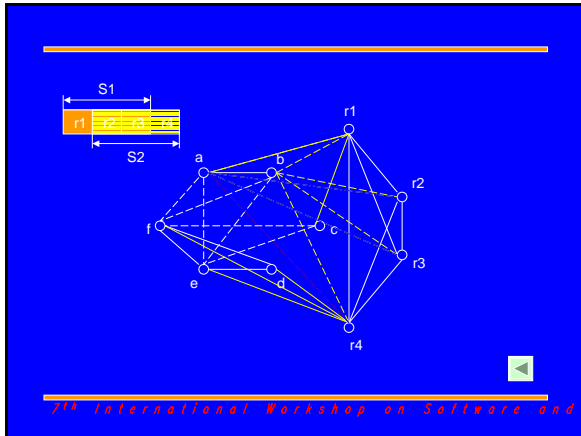




• Range assignment conflict:

- Strong conflict:  $u$  and  $r$  have strong conflict if  $u$  can never reside in  $S_i$ , where  $r$  in  $S_i$
- No conflict:  $u$  and  $r$  have no conflict if  $u$  can always reside in  $S_i$ , where  $r$  in  $S_j$
- Weak conflict:  $u$  and  $r$  have weak conflict if  $u$  can reside in  $S_j$ , where  $r$  not in  $S_j$





- Introduction
- The proposed architecture: LAR
- Sequential code generation for LAR
  - Annotated Conflict graph(ACG)
- The integrated approach
  - Annotated Worst-Case Conflict Graph(AWCCG)
- Experimental results
- Conclusions and future work

DFG_tuj	encoding			sequential			integrated		
	central	LAR	%	S	S_c	T(s)	S	S_c	T(s)
ar_filter_1,18	382	308	78.57	5	2	0.07	5	2	0.09
wdel_1,27	476	388	83.61	6	2	0.07	6	2	0.12
lact_2,20	714	588	82.35	9	5	inf	9	5	4.95
lact_4,11	714	588	82.35	9	6	inf	9	6	0.35
				9	7	0.26	9	7	0.37
loef_2,15	952	952	100	12	8	inf	12	8	1.43
				12	9	0.2	12	9	no
loef_4,11	952	952	100	12	9	inf	12	9	0.68
				8	4	inf	8	4	1.28
chen_2,15	680	560	82.35	8	5	0.24	8	5	4.28
				9	3	0.22	9	3	no
chen_4,8	680	560	82.35	9	4	0.19	9	4	0.24
				9	5	0.16	9	5	0.24

- Conclusions and future work
- Conclusions
  - New encoding style for reducing code size
  - No extra hardware, no extra move operations
  - Corresponding code generation techniques
  - ACG for range constraints
  - AWCCG solves phase coupling problem
- Future work
  - More versatile architectures
  - Combine with the operation assignment phase