

# Performance Analysis for Identification of (Sub)task-Level Parallelism in Java

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## **Performance Analysis:** need for embedded system program optimisation





## Outline

## >Introduction

□Parallel Performance Analysis (PPA)

□ Pre-processing

□ Profiler

□ Post-processing

Results & Conclusions



# We do task-level parallelism extraction from object-oriented programs

### = high-level platform-independent transformations

### "high-level"

>looking at the high-level program structures e.g. classes, methods

### "platform-independent"

>positive effect for multiprocessor systems in general

thus, we have to: >identify dominant parts of the program
>extract task-level parallelism
>evaluate the transformation effect





# **Performance analysis requirements**

### program perspective:

□ **same environment** for original and transformed programs (to take equal measure for both)

## platform perspective:

# exposing the parallel behaviour (to evaluate the optimisation effects)

### designer's perspective:

 as fast as possible (minimal run-time overhead)
 running on any platform (most preferably on my computer)
 easy to use





# **Concept of virtual time**





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# **Parallel Performance Analysis**

### □pre-processing

user-controlled instrumentationprogram transformation for profiler

## □parallel profiler

□implements the run-time support for the concept of virtual time

- ≻executing program
- >simulating parallel behaviour

## □post-processing

critical-path analysisfeedback for the parallelism extraction





# **Pre-processing enables selective profiling based on user interest**

user indicates the important parts of the program

- 1. top-level methods and loops accumulating most of the computation
- 2. looking inside in more detail

### instrumentation:

inserts profiler-specific code to reflect user's interest

### profiling modes:

- □ full profile for n top levels in the call graph
- □ selective profile
- □ sub-graph (branch) profile
- □ cumulative vs. non-cumulative method profile





# Pre-processing adapts program for the parallel profiler

virtual time is based on passing time stamps between tasks of the program

- two possible situations:
- □task-creation
  - □ sending time stamp from parent to child
- □task-synchronisation
  - updating the stamp between synced peers

**solution** = transformation of the Java synchronisation primitives into profilerspecific one (binary semaphore)

reducing extentions to run-time systemreducing run-time overhead





# Parallel profiler implements run-time support for the virtual time concept

#### virtual time core

passing appropriate time stamps between cooperating threads

### interface

- providing control over the parallel profiler
- passing information between the program and the profiler



program control changing time information passing time information





# Profiler core – minimal functionality enabling parallel profiling

#### interpreter

- □ Java interpreter
- extended to enable operations
  on thread timers
- having configurable time unit for different processors

### thread timers

 single timer per thread
 storing proper time information in the proper thread timer

### semaphores

the only way to pass time information between different thread timers





## How it works ...

#### real execution sequence



### virtual execution sequence





# Extension interface enables full runtime control over the profiler

### general purpose counters

*init, inc/dec, set/rst, get* □typical usage: □per method timer □method call counter

### semaphore interface init(state), P(), V()

### profiler interface

reificationstatistical informationconfiguring profiling mode





# Example: simulating fixed number of processors via profiler interface



#### non-shared tid = unlimited no. of processors



#### notes:

❑th0 and th1 share the same thread counter❑threads are schedules by the JVM scheduler



# **Post-processing analysis indicates the** potential to improve

### critical-path analysis

the most critical part is the one where reduction in its execution time has the highest impact on overall execution time of the program



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### task balance

the ideal partitioning creates parallel sub-tasks with "balanced" execution time, i.e., their idleness is minimised



# **Experimental results**

	speedup	# th	idle [%]	T <sub>ins</sub> [s]	notes
MPEG player	~ 2.3	5	20	30	imperative, data-dominant, static
3D engine v1	~ 4.1	8	23	31	00D, modular, interactive
3D engine v2	~ 4.6	18	36	31	
javac v1	1.1 – 1.2	7 - 12	0	210	OOD, recursion, complex
javac v2	1.4 - 1.9	21 - 32	25 - 34	210	
javac v3	1.8 - 2.3	<b>21 - 32</b>	21 - 32	210	



# Conclusions

parallel performance analysis framework for task-level parallelism extraction

Concept of virtual time simulating parallel behaviour of multithreaded programs

Common execution environment for original and transformed programs

□run-time overhead < 3%

### Future:

□data-access analysis



## Thank you