

# Improving Offset Assignment through Simultaneous Variable Coalescing

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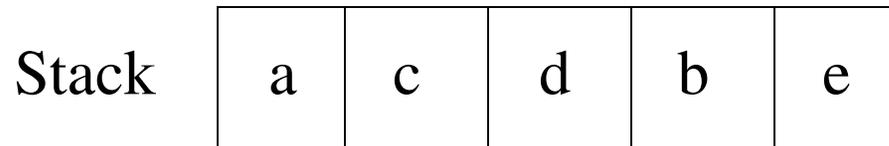
# Overview

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- Motivation
- Overview of Offset Assignment Problem
- Coalescing Simple Offset Assignment Technique
- Example of using CSOA Technique
- Experimental Results
- Conclusions and Future Work

# Motivation

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- Code size constraints
- Address computation is expensive
- ***Indirect addressing*** is suitable to embedded processors:
  - Implements ***fast address computation***
  - Enables the design of ***short instructions***
- ***Post-increment/decrement*** addressing modes

# Offset Assignment Problem

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- **Simple Offset Assignment Problem**: Almost all solutions are based in MWPC [Bartley'92, Liao et al'96, Leupers et al'96, Xiaotong'03, ...].
- **General Offset Assignment Problem**: Almost all solutions try to partition the variables and then resolves each partition as SOA problem [Liao et al'96, Leupers et al'96, Xiaotong'03, ...].

# Coalescing based Simple Offset Assignment

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- Based in the MWPC solution;
- At each step, chooses between:
  - (1) insert the maximal weight edge in the path; or
  - (2) coalesce two variables

# Coalescing based Simple Offset Assignment

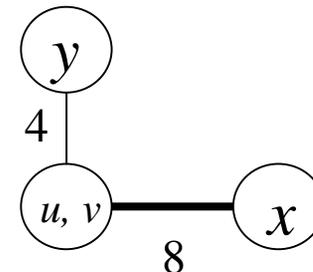
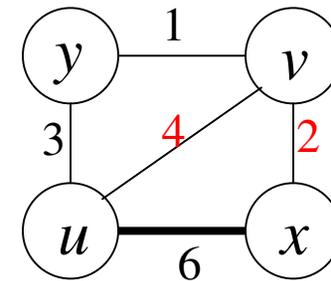
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- **Conditions to coalesce two variables  $u$  and  $v$ :**
  - $(u, v) \notin$  Interference Graph;
  - Coalesce  $u$  and  $v$  does not create a cycle, considering only the selected edges;
  - Coalesce  $u$  and  $v$  does not cause the new coalesced vertex to have degree greater than two, considering only the selected edges.

# Coalesce-based Simple Offset Assignment

## How to calculate the savings in offset cost when coalescing two variables $u$ and $v$ :

- For each  $x \in (Adj_{sel}(u) - Adj_{sel}(v))$ , add the weight of the edge  $(x, v)$  to the cost;
- For each  $x \in (Adj_{sel}(v) - Adj_{sel}(u))$ , add the weight of the edge  $(x, u)$  to the cost;
- Add the weight of the edge  $(u, v)$  to the cost, if this edge was not selected yet.



Offset Cost saved: 6

# Coalescing based Simple Offset Assignment

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**Algorithm 1** Coalescing-Based SOA

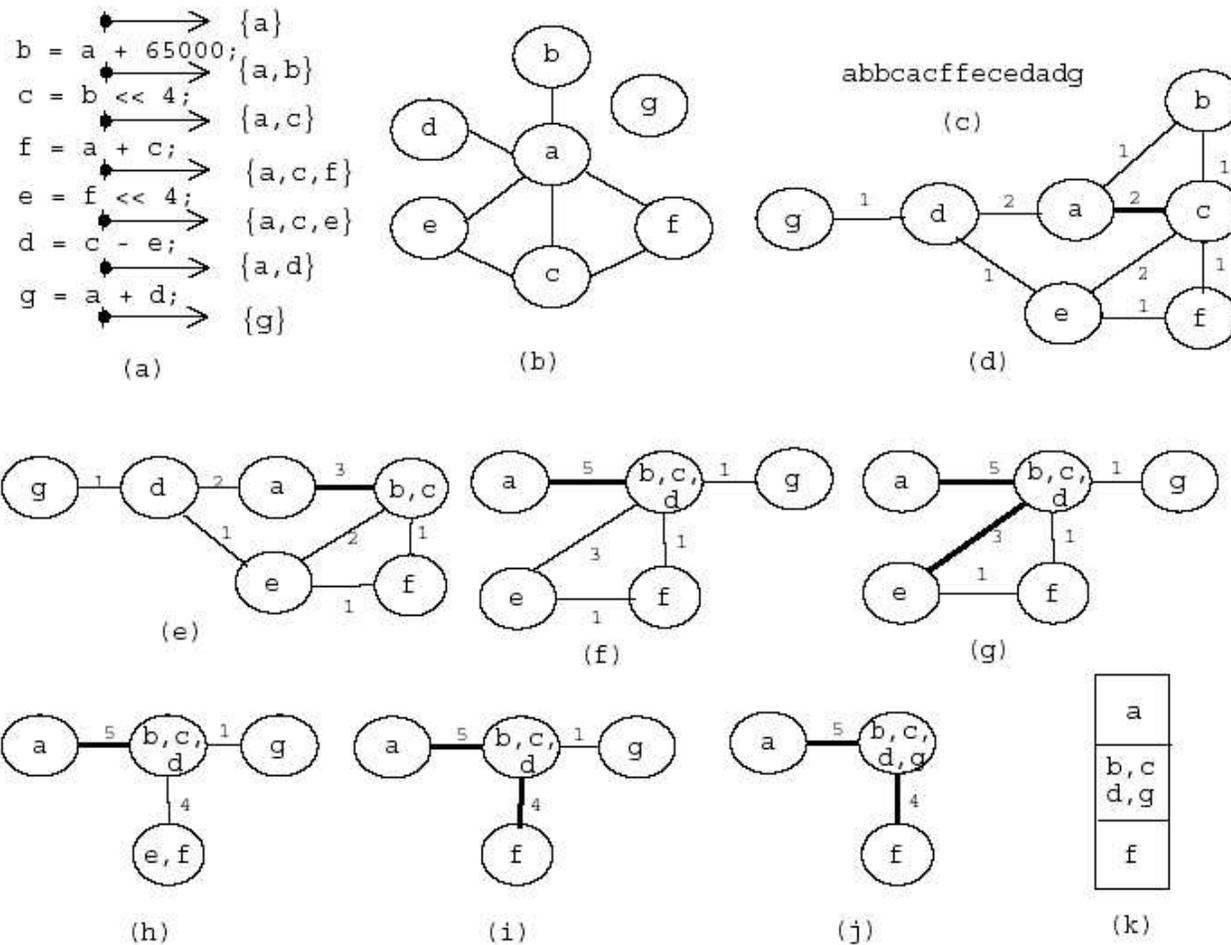
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(1)  $G_A(V_A, E_A) \leftarrow \text{BuildAccessGraph}(L_{AS});$ 
(2)  $L =$  sorted list of the  $E_A$ ;
(3)  $coal \leftarrow \text{false};$ 
(4)  $sel \leftarrow \text{false};$ 
(5) repeat
(6)    $rebuild \leftarrow \text{false};$ 
(7)    $(coal, u, v, csave) \leftarrow \text{FindCandidatePair}(G_I, u, v);$ 
(8)    $sel \leftarrow \text{FindEdgeValidNotSel}(L, e);$ 
(9)   if  $(coal \ \&\& \ sel)$ 
(10)    if  $(csave \geq w(e))$ 
(11)      $rebuild \leftarrow \text{true};$ 
(12)    else
(13)     mark  $e$  as selected;
(14)   else
(15)     if  $(coal)$ 
(16)       $rebuild \leftarrow \text{true};$ 
(17)     else if  $(sel)$ 
(18)      mark  $e$  as selected;
(19)   if  $(rebuild)$ 
(20)      $\text{RebuildAccessGraph}(G_A, u, v);$ 
(21)      $\text{RebuildInterferenceGraph}(G_I, u, v);$ 
(22)      $\text{RebuildL}(L);$ 
(23) until  $(!(coal \ || \ sel))$ 
(24) return  $\text{BuildOffset}(G_A);$ 
```

Input: the access sequence  $L_{AS}$ ,  
the interference graph  $G_I(V_I, E_I)$ .  
Output: the offset assignment.

# Coalescing based Simple Offset Assignment

- Example of algorithm execution**



# Experimental Results – Offset Costs (relative SOA-Liao costs [Liao et al'96])

- Implementation using LANCE and OffsetStone:
  - TB [Leupers et al'1996]
  - GA [Leupers et al'1998]
  - INC-TB [Leupers'2003]
  - SOA-Color [Ottoni et al'03]
  - CSOA [Ottoni et al'03]

Benchmarks	TB	GA	INC-TB	SOA-Color	CSOA
adpcm	89.1%	89.1%	89.1%	55.8%	45.6%
epic	96.8%	96.6%	96.6%	74.3%	50.2%
g721	96.2%	96.2%	96.2%	50.6%	27.9%
gsm	96.3%	96.3%	96.3%	26.6%	19.4%
jpeg	96.9%	96.7%	96.7%	52.6%	32.2%
mpeg2	97.3%	97.1%	97.2%	60.2%	34.3%
pegwit	91.1%	90.7%	90.7%	75.2%	38.8%
pgp	94.9%	94.8%	94.8%	55.0%	32.2%
rasta	98.6%	98.5%	98.5%	33.2%	21.1%
Average	95.2%	95.1%	95.1%	51.2%	32.1%

## Experimental Results – *Memory Savings*

Benchmarks	SOA-Color			CSOA		
	Code	Data	Code+Data	Code	Data	Code+Data
adpcm	89.9%	15.7%	71.5%	87.5%	27.8%	72.7%
epic	92.0%	11.6%	74.1%	84.6%	27.0%	71.8%
g721	90.7%	13.3%	68.0%	86.4%	25.1%	68.5%
gsm	94.3%	7.0%	71.9%	93.7%	21.8%	75.3%
jpeg	94.9%	14.6%	81.0%	92.7%	34.9%	82.7%
mpeg2	92.8%	12.5%	75.1%	88.0%	31.9%	75.7%
pegwit	97.4%	9.7%	78.2%	93.6%	35.3%	80.8%
pgp	99.6%	12.8%	98.3%	99.4%	31.6%	98.4%
rasta	84.4%	9.9%	71.2%	81.6%	26.1%	71.8%
Average	92.8%	11.6%	76.2%	89.6%	28.7%	77.1%

- Memory Savings relative to memory used by Liao's algorithm.

# Conclusions and Future Works

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## Conclusions:

- Contrary to what Liao suspected in his thesis (MIT'96), coalescing can improve OA.
- CSOA heuristic produces good results, dramatically reducing the OA cost.

## Future Work:

- CSOA can be used to resolve GOA problem.