## Communication-based Development of Systems Using Standard Programming Languages

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KPS'09

Communication-based Development with Standard Programming Languages

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#### Motivation

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# Communication-based specification



?

CODE

Program

message sequence charts sequence diagrams stream processing Java Haskell Erlang

. . .

#### Graphical Communication-based Specification

**Example** Transmission of a message *d* from a sender to a receiver via an unreliable transmission with acknowledgements



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#### Stream Processing

Streams record the sequence of messages transmitted via a communication line.

Finite streams: 
$$\mathcal{A}^{\star} = \underbrace{\{\langle \rangle\}}_{\text{empty}} \cup \underbrace{\mathcal{A} \times \mathcal{A}^{\star}}_{\text{non-empty}}$$

stream  $x \triangleleft X$ 

Infinite streams arise as limit points.

A stream function  $f : \mathcal{A}_1^* \times \ldots \times \mathcal{A}_m^* \to \mathcal{B}_1^* \times \ldots \times \mathcal{B}_n^*$  maps input streams to output streams:

streams



monotonicity Previous output messages can not be cancelled. continuity The behaviour on infinite input streams is approximated by the behaviour on finite streams.

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## Example Sender: Input / Output Behaviour

Interface send :  $\mathcal{D}^\star \times \mathcal{A}^\star \to \mathcal{D}^\star$ 

#### Behaviour

$$send(\langle \rangle, A) = \langle \rangle$$
  

$$send(d \triangleleft D, A) = d \triangleleft wait(d)(D, A)$$
  

$$wait(d)(D, \langle \rangle) = \langle \rangle$$
  

$$wait(d)(D, \oplus \triangleleft A) = send(D, A)$$
  

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#### **Eliminating** *wait*:

$$send(\langle \rangle, A) = \langle \rangle$$
  

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#### Communication-based Programming in Haskell

Laziness makes Haskell an interesting candidate for communication-based programming.

channels lazy lists

components list functions

composition network of mutually recursive equations naming the channels

! caution with patterns for lazy lists

## Example in Haskell

transmit ::  $[Int] \rightarrow [a] \rightarrow [Maybe a] \dots$ receive ::  $[Maybe a] \rightarrow ([a], [Bool]) \dots$ 

network oracle 
$$xs = ys$$
 xs send ds transmit cs ys  
where  $ds = send xs$  as  
 $cs = transmit oracle ds$   
 $(ys, as) = receive cs$ 

## Example in Haskell

 $\begin{array}{rrrr} \mbox{transmit} & :: & [\mbox{Int}] & -> & [\mbox{a}] & -> & [\mbox{Maybe a}] & \dots \\ \mbox{receive} & :: & [\mbox{Maybe a}] & -> & ([\mbox{a}], [\mbox{Bool}]) & \dots \end{array}$ 

network oracle 
$$xs = ys$$
  $xs$  send  $ds$  transmit  $cs$   $ys$   
where  $ds = send xs$  as  
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 $(ys, as) = receive cs$ 

works fine

#### Example in Haskell without Auxiliary wait Function

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network oracle 
$$xs = ys$$
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where  $ds = send xs$  as  
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 $(ys, as) = receive cs$   
**ERROR** (control stack overflow)

#### Communication-based Programming in Java

Threads make Java suitable for communication-based programming.

channels threads realizing FiFo queues components threads accessing input and output ports composition manual setting of channels as input and output ports support STREAMS! tool

State transition tables form the basis for the implementation of the components.

State transition table for the sender:

current state		input		next state		output
control	data	mess	ack	control	data	
send		d		wait	d	$\langle d \rangle$
wait	d		$\oplus$	send		
wait	d		$\ominus$	wait	d	$\langle d \rangle$

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#### Example in Java

A component must contain a method which implements the transitions of the state transition table:

```
public void processStep() throws ... {
  switch (cState) {
    case SENDING: if (!isEmpty(0)) {
         cState = CState.WAITING; }
         dState = ((Integer)get(0)).intValue();
         set(0,dState);
      break:
    case WAITING:
       break:
```

Communication-based specification



state	input	state	output

. . .

State-based specification

С	0	D	Е
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Program

message sequence charts sequence diagrams stream processing

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state transition machines	Java
state diagrams	Haskell
state charts	Erlang

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