



WESTFÄLISCHE
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Adding Weights to Dynamic Pushdown Networks

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Motivation

- Weighted Pushdown Systems [RSJM05]
 - Extend Pushdown Systems with Weights
 - Generalize interprocedural Dataflow Analysis
- Dynamic Pushdown Networks [BMOT05]
 - Model Network of parallel Processes
 - Dynamic Creation of Processes
- Weighted Dynamic Pushdown Networks
 - Combination of WPDS and DPN
 - Generalize interprocedural Dataflow Analysis of parallel Processes with dynamic Process Creation

Weighted Dynamic Pushdown System

- WDPN $\mathcal{W} = (\mathcal{M}, \mathcal{S}, f)$
 - DPN $\mathcal{M} = (P, \Gamma, \Delta)$
 - Configurations $(P\Gamma^*)^*$
 - Transition Rules $p\gamma \hookrightarrow c, c \in (P\Gamma^*)^*P\Gamma^*$
 - Semiring $\mathcal{S} = (D, \oplus, \odot, 0, 1)$
 - Weight Function $f : \Delta \rightarrow D$
- Interleaving Semantics
- Abstraction $\alpha(r_1 \dots r_n) = f(r_1) \odot \dots \odot f(r_n)$

Generalised Pushdown Predecessor Problem

$$\delta(C_1, C_2) = \alpha(\text{Paths}(C_1, C_2))$$

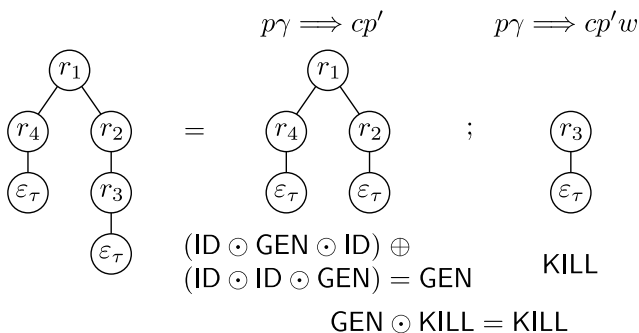
Executions of a DPN

$$\begin{array}{ll} C_1 = \{p\gamma_1\gamma_2\} & r_1 = p\gamma_1 \hookrightarrow p\gamma_3p\gamma_4 \quad \text{ID} \\ C_2 = \{p\gamma_6p\gamma_5\} & r_2 = p\gamma_4 \hookrightarrow p \quad \text{ID} \\ & r_3 = p\gamma_2 \hookrightarrow p\gamma_5 \quad \text{KILL} \\ & r_4 = p\gamma_3 \hookrightarrow p\gamma_6 \quad \text{GEN} \end{array}$$

$$p\gamma_1\gamma_2 \longrightarrow p\gamma_3p\gamma_4\gamma_2 \longrightarrow^* p\gamma_6p\gamma_5$$

$$\begin{array}{ll} \{r_1r_4r_2r_3, & (\text{ID} \odot \text{GEN} \odot \text{ID} \odot \text{KILL}) \oplus \\ r_1r_2r_4r_3, & (\text{ID} \odot \text{ID} \odot \text{GEN} \odot \text{KILL}) \oplus \\ r_1r_2r_3r_4\} & (\text{ID} \odot \text{ID} \odot \text{KILL} \odot \text{GEN}) = \text{GEN} \end{array}$$

Execution Hedges of a DPN



Reaching Paths and Hedges

$$\text{Paths}(C_1, C_2) = \psi(\text{Hedges}(C_1, C_2))$$

Branching Weighted Dynamic Pushdown System

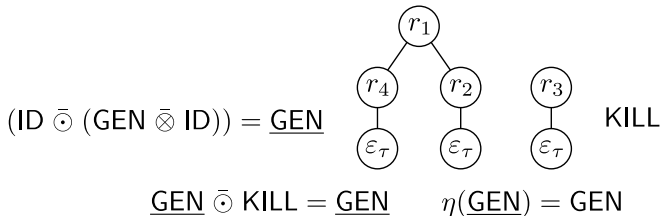
- BWDPN $\mathcal{B} = (\mathcal{M}, \mathcal{E}, \bar{f})$
 - DPN $\mathcal{M} = (P, \Gamma, \Delta)$
 - Extended Semiring $\mathcal{E} = (E, \bar{\oplus}, \bar{\odot}, \bar{\otimes}, \bar{0}, \bar{1})$
 - $(e_1 \bar{\otimes} e_2) \bar{\odot} e_3 = e_1 \bar{\otimes} (e_2 \bar{\odot} e_3)$
 - Weight Function $\bar{f} : \Delta \rightarrow E$
- Branching Semantics
- Abstraction $\beta(r(\tau_1 \dots \tau_n)) = \bar{f}(r) \bar{\odot} (\beta(\tau_1) \bar{\otimes} \dots \bar{\otimes} \beta(\tau_n))$

Branching Generalised Pushdown Predecessor Problem

$$\theta(C_1, C_2) = \beta(\text{Hedges}(C_1, C_2))$$

Connection

- Extension $(\mathcal{S}, \mathcal{E}, \iota, \eta)$
 - Embedding $\iota : D \rightarrow E$
 - Projection $\eta : E \rightarrow D$



Connection

$$\delta(C_1, C_2) = \eta(\theta(C_1, C_2)) \text{ für } \bar{f}(r) = \iota(f(r))$$

Applications/Future Work

- Applications
 - Shortest Path
 - Bitvector Analyses
 - KILL/GEN Analyses
- Future Work
 - Iterate Method
 - Synchronisation
 - Additional Weight Domains
 - Generalize Construction
 - Weights for Hedges

Bibliography



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