

Automatic Calculation of Coverage Profiles for Coverage-based Testing

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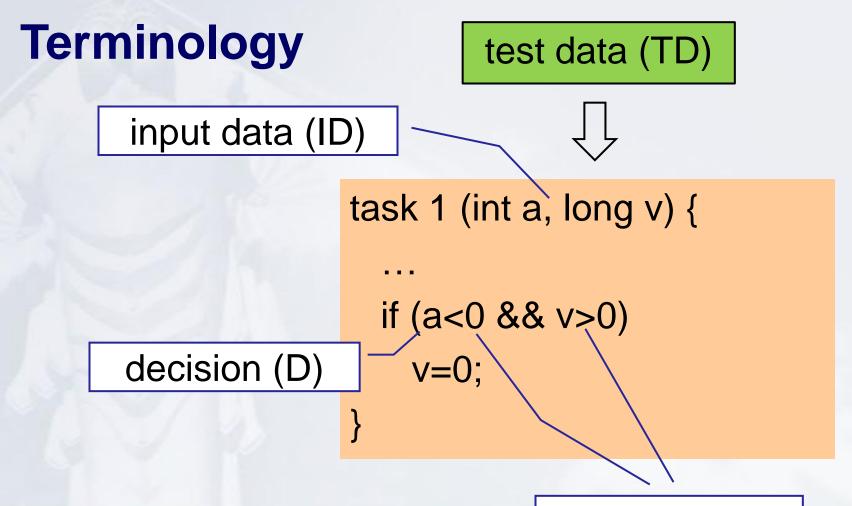
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Structural Code Coverage Metrics to Guide Software Testing

- Path Coverage: each path of the CFG is tested (infeasible in practice)
- State Coverage: each reachable state (variable values + program counter + HW states) is tested → most complex metrics (infeasible in practice)
- Statement Coverage: each statement is executed at least once
- Decision Coverage (aka. Branch Coverage): each edge of the CFG is executed at least once → tests conditional branches for both outcomes





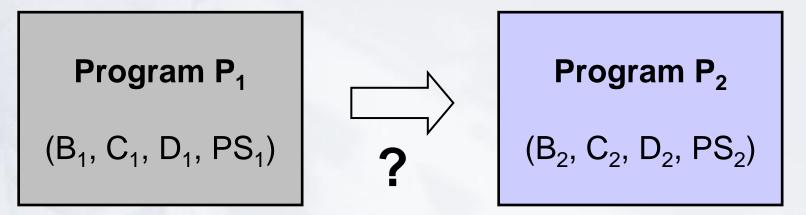








Question: Are Coverage Metrics Preserved by Code Transformations?



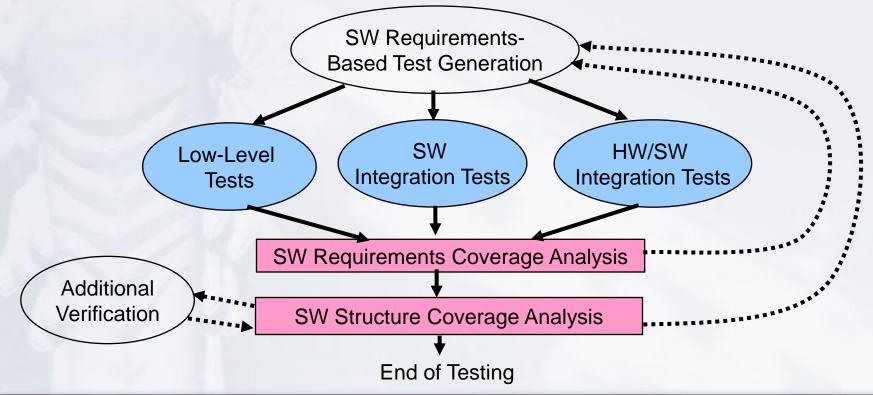
coverage(P_1 , TD) \equiv coverage(P_2 , TD)





RTCA/DO-178b (Civil Avionics)

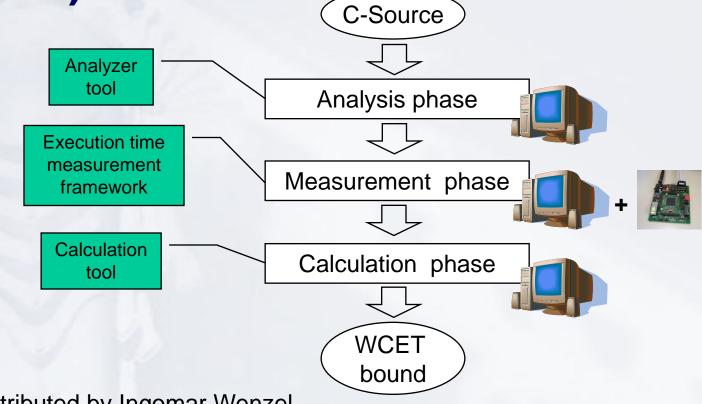
• Software testing process according to DO-178b:







Measurement-based Timing Analysis (MBTA)

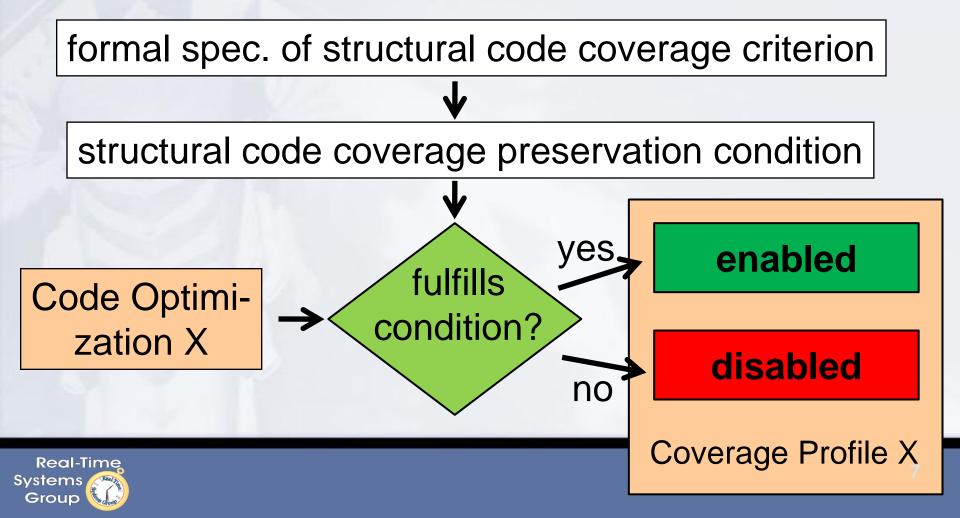


slide contributed by Ingomar Wenzel



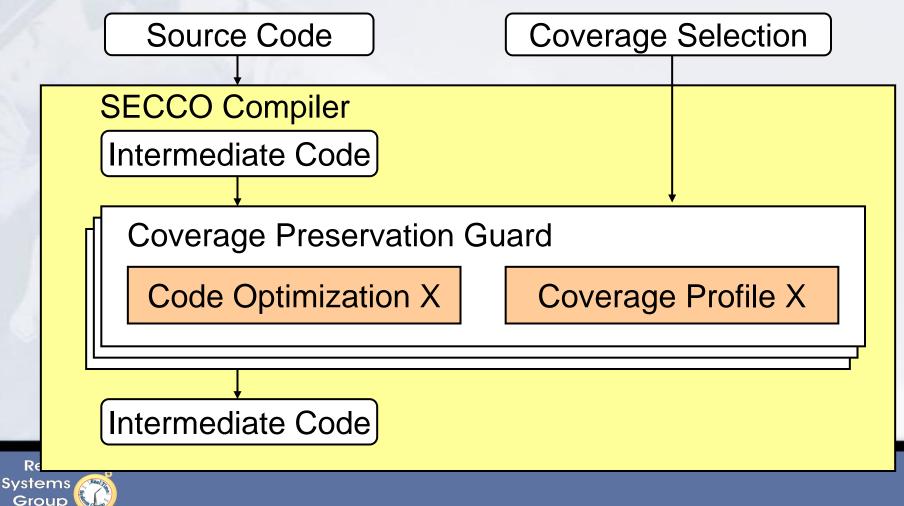


The SCCP/x Coverage Profiles





The SCCP/x Compilation Profile





Formal Notations

- Reachability Valuation: IV_R(x)
 ... set of valuations of input variables that trigger the execution of expression X.
- Reachability Valuation: IV_T(x), IV_F(x)
 ... set of valuations of input variables that trigger the execution of boolean expression X with a certain result of X:
 IV_T(x) ... x evaluates to TRUE
 IV_F(x) ... x evaluates to FALSE





Preservation of Decision Coverage

- Formal Definition of DC:
 ∀d∈D.(IV_T(d) ∩TD) ≠ φ ∧ (IV_F(d) ∩TD) ≠ φ
 D... set of decisions in the program
- DC Preservation Condition:
 ∀d₂∈D₂ ∃d₁∈D₁. touches_ID(d₁, IV_T(d₂)) ∧ ∃d₁∈D₁. touches_ID(d₁, IV_F(d₂))
 D₁... set of decisions of original program
 D₂... set of decisions of transformed program

touches_ID(x, ID) \Rightarrow $IV_T(x) \subseteq ID \lor IV_F(x) \subseteq ID$ (coverage of x implies that at least one test vector is element of input data set ID)

Analysis of Code Transformations

- A code transformation potentially disrupts a given structural code coverage if it:
 - changes the reachability of statements or conditions
 - adds new control-flow paths into the program





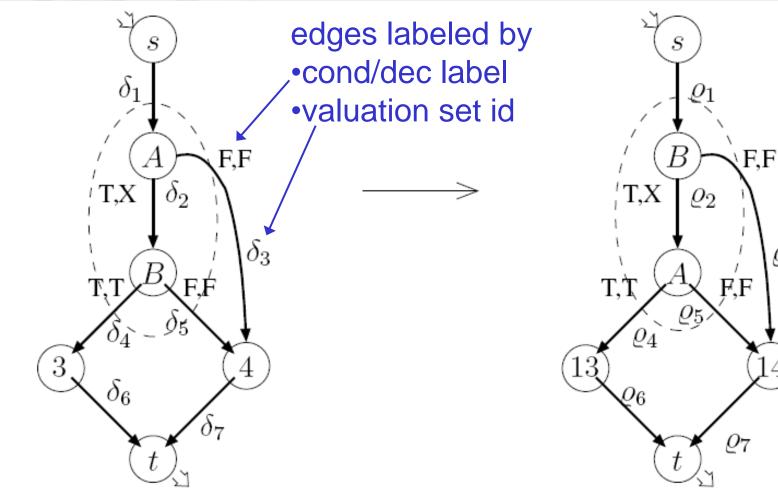
Example: Condition Reordering (with short-circuit)

1 if (A && B) 1 if (B && A)
2 thenBlock
3 else 3 else 3 else
4 elseBlock 4 elseBlock





Structural Transformation of Condition Reordering (with short-circuit)

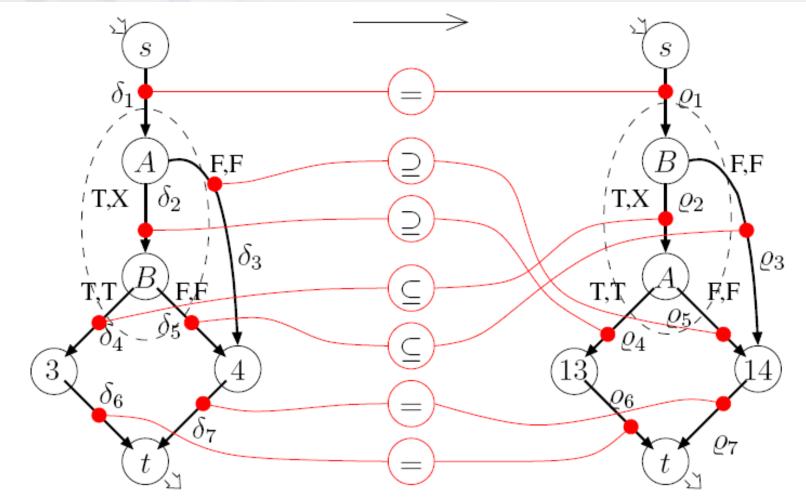


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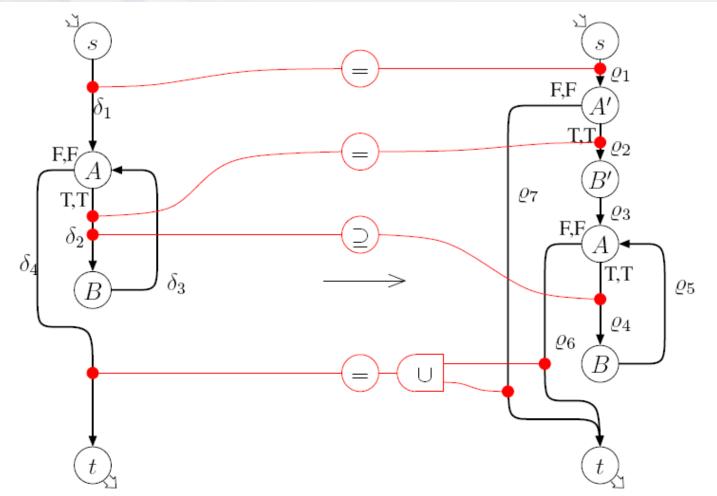


Transformation Relation for Condition Reordering (with short-circuit)





Transformation Relation for Loop Inversion



15



Results of Calculating Coverage Profiles

Transformation	SC	СС	DC
Condition reordering (without short-circuit)	\checkmark	•	\checkmark
Condition reordering (with short-circuit)	\checkmark	\checkmark	\checkmark
Loop peeling	•	·	•
Loop inversion	\checkmark	•	•





Future Work

- Modeling of additional (structural) code-coverage metrics (MCDC, scoped path coverage, ...)
- Calculation of coverage profiles for additional code transformations
- Modeling of code transformations close to the concrete implementation (reduce cognitive complexity of modeling)





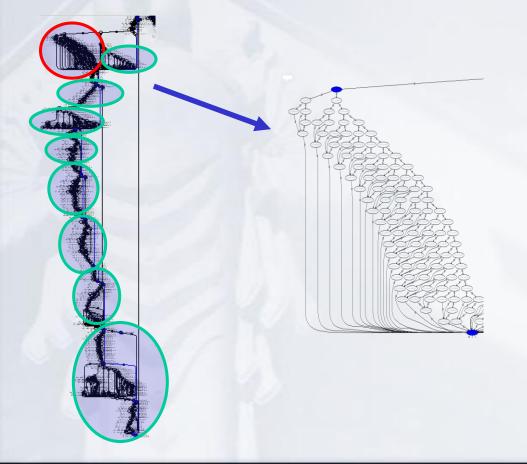
Further Reading

- [1] Raimund Kirner, SCCP/x A Compilation Profile to Support Testing and Verification of Optimized Code, In Proc. ACM Int. Conference on Compilers, Architecture, and Synthesis for Embedded Systems (CASES'07), Sep./Oct. 2007, pages 38-42, Salzburg, Austria.
- [2] Raimund Kirner and Susanne Kandl. Test coverage analysis and preservation for requirements-based testing of safety-critical systems. ERCIM News, (75):40–41, Oct. 2008.
- [3] Raimund Kirner. *Towards preserving model coverage and structural code coverage*. EURASIP Journal on Embedded Systems, Hindawi, 2009.





MBTA: CFG Partitioning



Example of generated code:

Entity	Value
#Paths	2.19e+32

(slide contributed by Ingomar Wenzel)



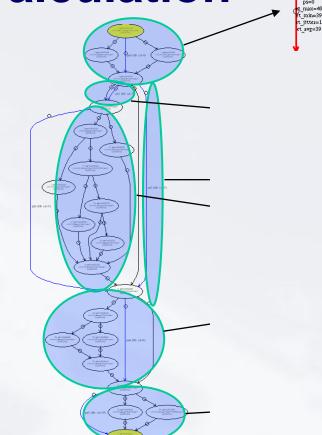


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MBTA: WCET Bound Calculation

- Program segment execution times are combined by integer linear programming (ILP) or longest path search
- Advantage: only feasible paths within PS contribute
- Deficiency: lack of global path information refinement possible



(slide contributed by Ingomar Wenzel)

