

Exercise 1 : (4 Points)

Prove by structural induction (on the inductive structure of arithmetic expressions) the Substitution Lemma 1.7.3 of Chapter 1.7:

Lemma 1.7.3 (*Substitution Lemma for $\llbracket \cdot \rrbracket_A$*)

$$\forall a, a' \in \mathbf{Aexpr}. \forall \sigma \in \Sigma. \llbracket a[a'/x] \rrbracket_A(\sigma) = \llbracket a \rrbracket_A(\sigma[\llbracket a' \rrbracket_A(\sigma)/x])$$

Exercise 2 : (4+4 Points)

Let $\sigma \in \Sigma$ be a state with $\sigma(x) = \mathbf{13}$ and $\sigma(y) = \mathbf{5}$. Prove using the

1. structurally operational
2. natural

semantics of WHILE that the program

$z := 0; \mathbf{while} \ y \leq x \ \mathbf{do} \ z := z + 1; \ x := x - y \ \mathbf{od}$

applied to σ terminates regularly in state σ' with $\sigma' = \sigma[\mathbf{2}/z][\mathbf{5}/y][\mathbf{3}/x]$.

Exercise 3 : (4 Points)

Let $\pi_1, \pi_2 \in \mathbf{Prg}$ be two WHILE-programs, and let $\sigma, \sigma' \in \Sigma$ be two states. Investigate the validity of the below implication (proof or counterexample):

$$\langle \pi_1; \pi_2, \sigma \rangle \Rightarrow^* \langle \pi_2, \sigma' \rangle \Rightarrow \exists k \in \mathbb{N}_0. \langle \pi_1, \sigma \rangle \Rightarrow^k \sigma'$$

Exercise 4 : (4+4 Points)

Extend the programming language WHILE by the loop statement

$\mathbf{repeat} \ \pi \ \mathbf{until} \ b \ \mathbf{end}$

Provide a

1. SOS rule [rep_{sos}]
2. NS rule [rep_{ns}]

such that the statement gets its “usual” semantics without relying on or exploiting the existence of the while statement in WHILE for defining the two rules.

Submission: Wednesday, 20 March 2019, before the lecture.