
Today's Topics

- Part I: Type Checking/Type Inference
A necessity of compilers and interpreters
- Part II: Parallelism in Functional Programming Languages
A hot research topic
- Part III: The Story of Haskell
Behind the scenes of Haskell (and Functional Programming)

Part I: Type Checking/Type Inference

...of central importance for implementing functional programming languages like Haskell (as well as languages of other paradigms)

We distinguish...

- monomorphic and
- polymorphic

type checking/type inference.

Type Checking/Type Inference

Important, too, in this context...

- Overloading (as a special case of polymorphism (aka ad hoc polymorphism))
- Type classes (such as `Num`, `Eq`, etc.)

Reference

The following presentation is based on...

- Chapter 13
Simon Thompson. *Haskell – The Craft of Functional Programming*, Addison-Wesley, 2nd edition, 1999.
- Chapter 5
Martin Erwig. *Grundlagen funktionaler Programmierung*. Oldenbourg Verlag, 1999 (In German).

Type-Based Programming Languages

We distinguish...

- Programming languages with...
 - *weak* (checked at run-time)
 - *strong* (checked at compile-time)
- typing

Benefits

...of using typed programming languages

- *More reliable code*:
 - ...many programming flaws can be detected at compile-time; type correctness is a proof of correctness on the abstraction level of types
- *More efficient code*: ...no type-checks required at run-time
- *More effective program development*:
 - ...Type information is additional program documentation
 - ...the *understanding*, *advancing* and *maintaining* of programs gets simpler, e.g. the search for pre-defined library functions (“is there a library function, which removes duplicates from a list, i.e., applied to the list [2,3,2,1,3,4] yields the result [2,3,1,4]? Search for a function with type (Eq a) => [a] -> [a]”.)

What’s it all about? 1(4)

Haskell expressions all have a defined type.

This type can be...

- monomorphic
- polymorphic
 - restricted by type class constraints

...and *explicitly* be given as in the below examples:

```
'w'  :: Char           -- monomorphic
flip :: (a -> b -> c) -> (b -> a -> c) -- polymorphic
elem :: Eq a => a -> [a] -> Bool      -- polymorphic with
                                   type class constraint
```

What’s it all about? 2(4)

Types (of expressions) can automatically be inferred by Haskell compilers and interpreters as in the below example:

```
magicType = let
    pair x y z = z x y
    f y = pair y y
    g y = f (f y)
    h y = g (g y)
  in h (\x->x)
```

What's it all about? 3(4)

The call of `:t magicType` in Hugs yields:

```
Main> :t magicType
magicType ::
  (((((((a -> a) -> (a -> a) -> b) -> b) ->
  (((a -> a) -> (a -> a) -> b) -> b) -> c) -> c) ->
  (((((a -> a) -> (a -> a) -> b) -> b) ->
  (((a -> a) -> (a -> a) -> b) -> b) -> c) -> c) -> d) -> d) ->
  (((((((a -> a) -> (a -> a) -> b) -> b) -> (((a -> a) ->
  (a -> a) -> b) -> b) -> c) -> c) -> (((((a -> a) ->
  (a -> a) -> b) -> b) -> (((a -> a) ->
  (a -> a) -> b) -> b) -> c) -> c) -> d) -> d) -> e) -> e
```

Quite a complex type, isn't it?

What's it all about? 4(4)

Central question...

- How does Hugs succeed to automatically infer this type?

The systematic investigation leads us to notions such as...

- *Type analysis/-checking*
- *Type systems* and
- *Type inference*

First, an informal approach, driven by examples...

Monomorphic Type Checking

We first consider a simplified situation without...

- polymorphism

Characteristic for this situation is...

- An expression has
 - either a unique single type
 - or no type at all (not well-typed)

Convention:

The polymorphism of polymorphic or overloaded functions is explicitly resolved (by indexing) as in...

```
+_ {Int} :: Int -> Int -> Int
length_ {Char} :: [Char] -> Int
```

Type Analysis

Central idea...

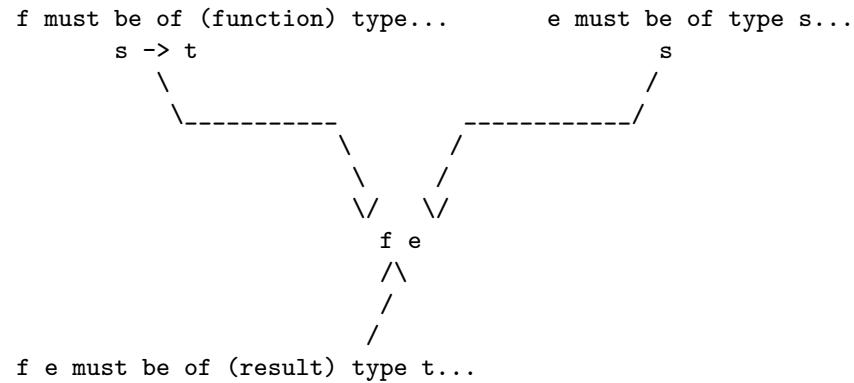
- Evaluate the application context an expression is embedded in

In the following a couple of examples for illustration...

Type Checking Expressions 1(3)

(A) Function Applications...

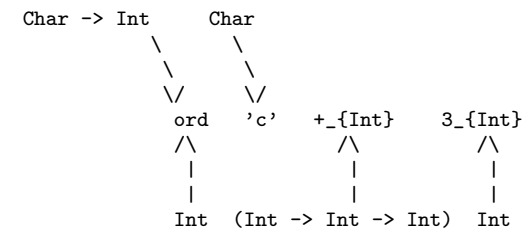
The context of function applications allows conclusions about the types involved...



Type Checking Expressions 2(3)

(B) Other expressions...

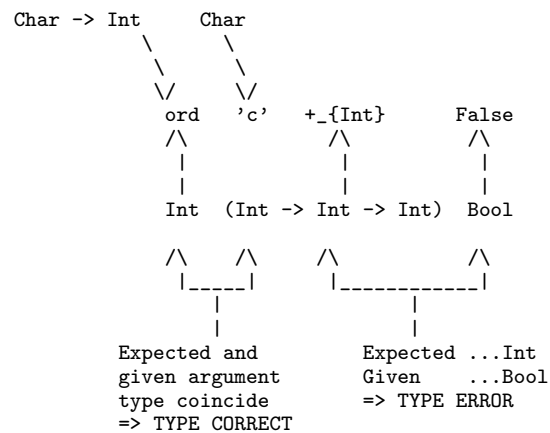
First, a correctly typed expression (ord 'c' +_Int 3_Int)...



Type Checking Expressions 3(3)

(B) Other expressions...

Second, an incorrectly typed expression (ord 'c' +_Int False)...



Type Checking Function Definitions

```
f :: t1 -> t2 -> ... -> tk -> t
f a1 a2 ... ak
  | b1 = e1
  | b2 = e2
  ...
  | bk = ek
```

In type checking monomorphic function definitions, three things have to be checked...

- Each guard b_i must be of type `Bool`
- The result value of each expression e_i must be of type t
- The pattern of each argument p_i must be consistent with the type of that formal parameter, i.e., with type t_i .

Consistency of Patterns 1(2)

Informally:

A pattern is *consistent* with a type, if it *matches* (some of) the values of this type.

In detail:

- A variable is consistent with any type
- A literal is consistent with its type
- A pattern $p:q$ is consistent with type $[t]$, if p is consistent with type τ and q is consistent with type $[t]$
- ...

Consistency of Patterns 2(2)

Example:

- $(42:xs)$...is consistent with $[Int]$
- $(x:xs)$...is consistent with any type of lists

Polymorphic Type Checking 1(2)

Characteristic for the situation of polymorphic type checking...

- An expression can
 - have several types (while being well-typed)

Central for algorithmically solving the polymorphic type checking problem is...

- *Constraint satisfaction*

which itself is based on

- *Unification*

Polymorphic Type Checking 2(2)

Consider...

```
length :: [a] -> Int
```

Informal interpretation of the type

```
length :: [a] -> Int
```

An abbreviation of (the set of types)

```
length :: [t] -> Int
```

where t is some arbitrary monomorphic type; in total, it is thus an abbreviation of the set of types

```
[Int] -> Int
[(Bool,Char)] -> Int
etc.
```

Polymorphic Type Checking – Continuing the Example

In the example of...

```
length ['c', 'a']
```

we can infer from the calling context the type...

```
length :: [Char] -> Int
```

Observation

The preceding examples allow us to conclude:

The application contexts of expressions impose...

- different constraints on the type of the expression.

This way type checking boils down to the problem, if...

- types can be determined for the various expressions such that all constraints are met.

Further Examples / Example 1

Consider:

```
f (x,y) = (x , ['a' .. y])
```

Observation:

- `f` ...expects pairs as arguments, where
 - 1st component: no imposed constraints
 - 2nd component: `y` must be of type `Char` because of being used in the range of an enumerated type `['a' .. y]`

Hence, we can infer the type of `f`:

```
f :: (a , Char) -> (a , [Char])
```

Example 2

Consider:

```
g (m,zs) = m + length zs
```

Observation:

- `g` ...expects pairs as arguments, where
 - 1st component: `m` must be of a numerical type because of being used as an operand of `+`
 - 2nd component: `zs` must be of type `[b]` because of being used as an argument of the function `length` with `length :: [b] -> Int`

Hence, we can infer the type of `g`:

```
g :: (Int, [b]) -> Int
```

Example 3 1(2)

Consider the composition of the two preceding examples:

$g . f$

Observation:

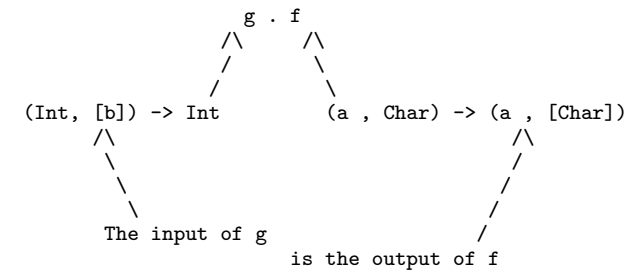
- $g . f$...in a composition $g . f$, the return value of f is passed as to g

In this example, this means...

- Result type of f : $(a, [\text{Char}])$
- Argument type of g : $(\text{Int}, [b])$
- Required: types for a and b which satisfy the above two constraints

Example 3 2(2)

Illustration:



Hence, we can infer the type of $g . f$:

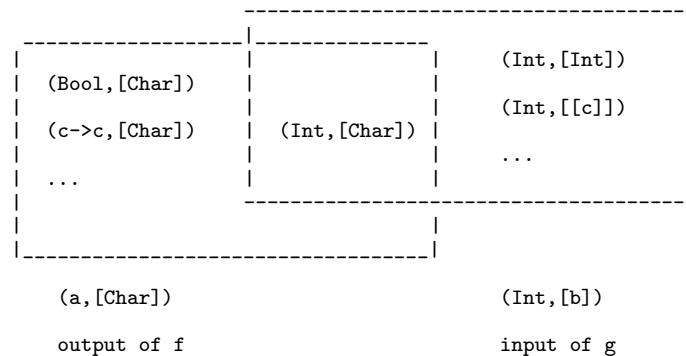
$g . f :: (\text{Int}, [\text{Char}]) \rightarrow \text{Int}$

Unification in the case of the Example

Crucial for drawing the conclusion in the previous example is...

- Unification

Illustration:



Unification

Introducing terminology. We say...

- *Instance* of a type ...given by replacing a type variable by a (concrete) type expression
- *Common instance* of two type expressions ...if the instance is an instance of both type expressions

Unification problem...

- Search for the *most general common (type) instance*

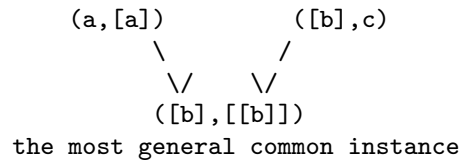
In the previous example:

- The most general common type instance of $(a, [\text{Char}])$ and $(\text{Int}, [b])$ is the single type $(\text{Int}, [\text{Char}])$.

More on Unification 1(3)

In general, unification does not lead to unique types...

Example:

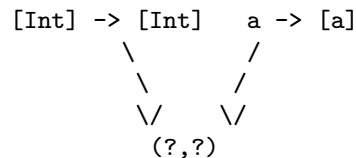


Observation:

- Constraint $(a, [a])$ requires: the second component is a list of elements of the type of the elements of the first component
- Constraint $([b], c)$ requires: the first component is of some list type
- Together, this implies: the most general common type instance of $(a, [a])$ and $([b], c)$ is $([b], [[b]])$

More on Unification 3(3)

Illustrating failure of unification in more detail...



Unification of the...

- Argument type requires ...a must be of type `[Int]`
- Result type requires ...a must be of type `Int`
- Together ...unification fails (inconsistent constraints)

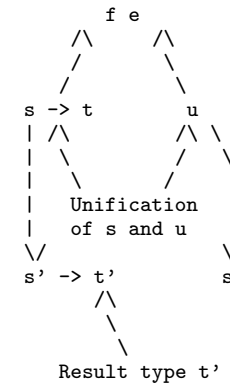
More on Unification 2(3)

Note:

- Instance \neq Unifier
 $([Bool], [[Bool]]), ([c], [[c]])$ are...
 - instances of $([b], [[b]])$
 - but no unifier: $([b], [[b]])$ is not an instance of either of them
- Unification can fail: $[Int] \rightarrow [Int]$ and $a \rightarrow [a]$ aren't unifiable

Type Checking Expressions

Here... polymorphic function application



Observation:

- s and u need not be equal; they only need to be *unifiable*

Typ Checking and Type Classes

Consider...

```
member []      y = False
member (x:xs) y = (x==y) || member xs y
```

In the above example, the usage of (==) forces...

```
member :: Eq a => [a] -> a -> Bool
```

Type Checking and Overloading

Consider the application of function members to e:

```
member e
```

with

```
e      :: Ord b => [[b]]
```

Without further context information unification yields...

```
member :: [[b]] -> [b] -> Bool
```

Thus, we receive...

```
member e :: [b] -> Bool
```

More carefully, we have to take the contexts into account, too!

Context Analysis

```
(Eq [b] , Ord b)
```

Analysing and simplifying the contexts yields...

- Context constraints refer to type variables

```
instance Eq a => Eq [a] where...
...this yields (Eq b, Ord b)
```
- Repeat until no more instances apply
- Simplification of the context by means of the information given by class yields...

```
class Eq a => Ord a where...
```
- Hence ...Ord b
- Summing up ...

```
member e :: Ord b => [b] -> Bool
```

Summary – A Three-Stage Process

The three-stage process consists of

- Unification
- Analysis (with instances)
- Simplification

...is a typical pattern of the context-aware type analysis in Haskell.

Type Systems and Type Inference

Informally...

- *Type systems* are...
 - logical systems, which allow us to formalize statements of the form “exp is of type t” and to prove them by means of the axioms and rules of the type system
 - *Type inference* denotes...
 - the process to automatically derive the type of an expression by means of the axioms and rules of a type system
- Key words:* Type inference algorithm, Unification

A Typical Part of a Type Grammar

...generates the language of types

$$\begin{aligned} \tau ::= & \text{Int|Float|Char|Bool} && \text{(simple type)} \\ & | \alpha && \text{(type variable)} \\ & | \tau \rightarrow \tau && \text{(function type)} \\ \\ \sigma ::= & \tau && \text{(type)} \\ & | \forall \alpha. \sigma && \text{(type binding)} \end{aligned}$$

We say:

- τ ...a type
- σ ...a type schema

A Typical Part of a Type System 1(2)

...associates with each (typable) expression of the language a type of the type language

$$\begin{aligned} \text{VAR} & \frac{}{\Gamma \vdash \text{var} : \Gamma(\text{var})} \\ \text{CON} & \frac{}{\Gamma \vdash \text{con} : \Gamma(\text{con})} \\ \text{COND} & \frac{\Gamma \vdash \text{exp} : \text{Bool} \quad \Gamma \vdash \text{exp}_1 : \tau \quad \Gamma \vdash \text{exp}_2 : \tau}{\Gamma \vdash \text{if } \text{exp} \text{ then } \text{exp}_1 \text{ else } \text{exp}_2 : \tau} \\ \text{APP} & \frac{\Gamma \vdash \text{exp} : \tau' \rightarrow \tau \quad \Gamma \vdash \text{exp}' : \tau'}{\Gamma \vdash \text{exp } \text{exp}' : \tau} \\ \text{ABS} & \frac{\Gamma[\text{var} \mapsto \tau'] \vdash \text{exp} : \tau}{\Gamma \vdash /x \rightarrow \text{exp} : \tau' \rightarrow \tau} \\ & \dots \end{aligned}$$

where Γ is a so-called *type assumption*.

Typical Part of a Type System 2(2)

Type assumptions are...

- partial functions, which map variables to type schemas

Here, $\Gamma[\text{var}_1 \mapsto \tau_1, \dots, \text{var}_n \mapsto \tau_n]$ is the function, which yields the type τ_i for each var_i and is as Γ otherwise.

The Schematic Unification Algorithm

$$\begin{aligned} \mathcal{U}(\alpha, \alpha) &= [] \\ \mathcal{U}(\alpha, \tau) &= \begin{cases} [\tau/\alpha] & \text{if } \alpha \notin \tau \\ \text{error} & \text{otherwise} \end{cases} \\ \mathcal{U}(\tau, \alpha) &= \mathcal{U}(\alpha, \tau) \\ \mathcal{U}(\tau_1 \rightarrow \tau_2, \tau_3 \rightarrow \tau_4) &= \mathcal{U}(\mathcal{U}\tau_2, \mathcal{U}\tau_4)\mathcal{U} \\ &\quad \text{where } \mathcal{U} = \mathcal{U}(\tau_1, \tau_3) \\ \mathcal{U}(\tau, \tau') &= \begin{cases} [] & \text{if } \tau = \tau' \\ \text{error} & \text{otherwise} \end{cases} \end{aligned}$$

Remarks:

- \mathcal{U} ... (most general) unifier (essentially a substitution)
- Application of equations sequentially from top to bottom

Example on Unification/Most General Unification

Consider... $a \rightarrow (\text{Bool}, c)$ and $\text{Int} \rightarrow b$

- *Unifier* ... Substitution $[\text{Int}/a, \text{Float}/c, (\text{Bool}, \text{Float})/b]$
- *Most general unifier* ... Substitution $[\text{Int}/a, (\text{Bool}, c)/b]$

Example on the Unification Algorithm

Task

...Unifying the type expressions $a \rightarrow c$ and $b \rightarrow \text{Int} \rightarrow a$

Solution

$$\begin{aligned} \mathcal{U}(a \rightarrow c, b \rightarrow \text{Int} \rightarrow a) \\ (\text{mit } \mathcal{U} = \mathcal{U}(a, b) = [b/a]) &= \mathcal{U}(\mathcal{U}c, \mathcal{U}(\text{Int} \rightarrow a))\mathcal{U} \\ &= \mathcal{U}(c, \text{Int} \rightarrow b)[b/a] \\ &= [\text{Int} \rightarrow b/c][b/a] \\ &= [\text{Int} \rightarrow b/c, b/a] \end{aligned}$$

In total ... $b \rightarrow \text{Int} \rightarrow b$

Essence of Automatic Type Inference Algorithms

...syntax-directed application of the rules of the type inference systems

The key...

- Modifying the type inference system such that there is always only a single rule applicable

Further Reading 1(3)

Types and type systems, type inference...

- For functional languages in general
 - Anthony J. Field, Peter G. Robinson. *Functional Programming*. Addison-Wesley, 1988 (Chapter 7)
- Haskell-specific
 - Simon Peyton Jones, John Hughes. *Report on the Programming Language Haskell 98*.
<http://www.haskell.org/report/>

Further Reading 2(3)

- Overview
 - J. C. Mitchell. *Type Systems for Programming Languages*. In J. van Leeuwen (Hrsg.). *Handbook of Theoretical Computer Science, Vol. B: Formal Methods and Semantics*. Elsevier Science Publishers, 367-458, 1990.
- Foundations of polymorphic type systems
 - Robin Milner. *A Theory of Type Polymorphism in Programming*. *Journal of Computer and System Sciences* 17, 248-375, 1978.
 - L. Damas, Robin Milner. *Principal Type Schemes for Functional Programming Languages*. In *Conference Record of the 9th Annual ACM Symposium on Principles of Programming Languages (POPL'82)*, 207-218, 1982

Further Reading 3(3)

- Unification algorithm
 - J. A. Robinson. *A Machine-Oriented Logic Based on the Resolution Principle*. *Journal of the ACM* 12(1), 23-42, 1965.
- Type systems and type inference
 - Luca Cardelli. *Basic Polymorphic Type Checking*. *Science of Computer Programming* 8, 147-172, 1987.

Part II: Parallelism in Functional Programming Languages

Parallelism

- Implicit
- Explicit
- Skeletons

Reference

The following presentation is based on...

- Chapter 21
Peter Pepper, Petra Hofstedt. *Funktionale Programmierung*, Springer, 2006. (In German).

Related and relevant in this context...

- Murray Cole. *Algorithmic Skeletons: Structured Management of Parallel Computation*, The MIT Press, 1989.
- Philip W. Trinder, Kevin Hammond, Hans-Wolfgang Loidl, Simon L. Peyton Jones. *Algorithms + Strategy = Parallelism*. *Journal of Functional Programming*, 8(1):23-60, 1998.
- Philip W. Trinder, Hans-Wolfgang Loidl, Robert F. Poynton. *Parallel and Distributed Haskell*. *Journal of Functional Programming*, 12(4&5):469-510, 2002.

Parallelism in Functional Languages

In particular...

- Implicit (expression) parallelism
- Explicit parallelism
- Algorithmic skeletons

Parallelism in Imperative Languages

In particular...

- Data-parallel Languages (e.g. High Performance Fortran)
- Libraries (PVM, MPI) \rightsquigarrow *Message Passing Model* (C, C++, Fortran)

Implicit Parallelism

...aka *expression parallelism*

Consider the functional expression of the form $f(e_1, \dots, e_n)$:

Note:

- Arguments (and functions) can be evaluated in parallel.
- Advantages: Parallelism *for free!* No effort for the programmer.
- Disadvantages: Results often unsatisfying. E.g. granularity, load distribution, etc. not taken into account.

Thus:

- Easy to detect parallelism (i.e., for the compiler), but hard to fully exploit.

Explicit Parallelism

By...

- Introducing meta-statements (e.g. to control the data and load distribution, communication)
- Advantages: Often superior results by explicit hands-on control of the programmer
- Disadvantages: High programming effort, loss of functional elegance

Algorithmic Skeletons

A compromise between...

- *explicit imperative* parallel programming
- *implicit functional* expression parallelism

In the following

- Massively parallel systems
- Algorithmic skeletons

Massively Parallel Systems

...characterized by

- large number of processors with
 - local memory
 - communication by message exchange
- MIMD-Parallel Processor Architecture (*Multiple Instruction/Multiple Data*)
- Here we restrict ourselves to: SPMD-Programming Style (*Single Program/Multiple Data*)

Algorithmic Skeletons

Algorithmic Skeletons...

- represent typical patterns for parallelization (*Farm, Map, Reduce, Branch&Bound, Divide&Conquer,...*)
- are easy to instantiate for the programmer
- allow parallel programming at a high level of abstraction

Realization of Algorithmic Skeletons

...in functional languages

- by special higher-order functions
- with parallel implementation
- embedded in sequential languages

Thus

- Hiding of parallel implementation details in the skeleton
- Elegance and (parallel) efficiency for special application patterns

Example: Parallel Map on Distributed List

Consider the higher-order function `map` on lists...

```
map :: (a -> b) -> [a] -> [b]
map _ [] = []
map f (x:xs) = (f x) : (map f xs)
```

Observation

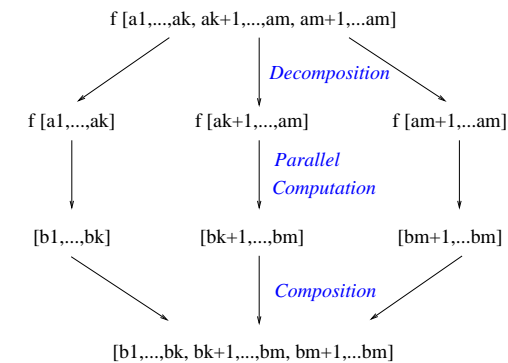
- Applying `f` to a list element does not depend on other list elements

Obvious

- Dividing the list into sublists followed by *parallel* application of `map` to the sublists (parallelization pattern *Farm*)

Parallel Map on Distributed Lists

For illustration...



On the Implementation

Implementing the parallel map function requires...

- special data structures, which take into account the aspect of distribution (ordinary lists are inefficient for this purpose)

Skeletons on distributed data structures

- so-called *data-parallel skeletons*

Note the difference:

- *Data-parallelism*: Assumes an *a priori* distribution of data on different processors
- *Task-parallelism*: Processes and data to be distributed are not known *a priori*, hence dynamically generated

Programming of a Parallel Application

...using algorithmic skeletons

- Recognizing problem-inherent parallelism
- Selecting an adequate data distribution (granularity)
- Selecting a suitable skeleton from a library
- Problem-specific instantiation of the skeleton(s)

Remark:

- Some languages (e.g. Eden) support the implementation of skeletons (in addition to those which might be provided by a library)

Data Distribution on Processors

...is

- crucial for
 - structure of the complete algorithm
 - efficiency

The hardness of the distribution problems depends on...

- Independence of all data elements (like in the map-example): Distribution is easy
- Independence of subsets of data elements
- Complex dependences of data elements: Adequate distribution is challenging

An auxiliary means

- So-called *covers* (investigated by various authors)

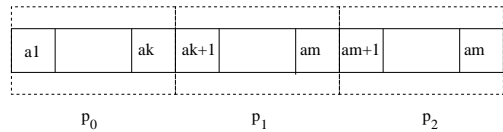
Covers

...describe the

- decomposition and communication pattern of a data structure

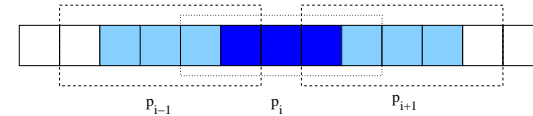
Example: Simple List Cover

Distributing a list on 3 processors p_0 , p_1 , and p_2 :



Peter Pepper, Petra Hofstedt. Funktionale Programmierung.
Springer, 2006, S. 446.

Example: List Cover with Overlapping Elements



Peter Pepper, Petra Hofstedt. Funktionale Programmierung.
Springer, 2006, S. 446.

General Cover Structure

```
Cover = {
  Type S a      -- Whole object
  C b          -- Cover
  U c          -- Local sub-objects

  split :: S a -> C (U a)  -- Decomposing the original object
  glue  :: C (U a) -> S a  -- Composing the original object
}
```

It is required:

```
glue . split = id
```

Note: No (valid) Haskell

Realization in a Programming Language

...implementing covers requires support for

- the specification of covers
- the programming of algorithmic skeletons on covers
- the provision of often used skeletons in libraries

...is

- currently a hot research topic in functional programming

Last but not least

Implementing skeletons...

- by message passing via skeleton hierarchies

Further Reading

- Hans-Werner Loidl et al. *Comparing Parallel Functional Languages: Programming and Performance*. Higher-Order and Symbolic Computation, 16(3):203-251, 2003.

Part III: The Story of Haskell

16 Years of Haskell: A Retrospective on the occasion of its
15th Anniversary

by

Simon Peyton Jones

Wearing the Hair Shirt: A Retrospective on Haskell

<http://research.microsoft.com/users/simonpj/papers/haskell-retrospective/>

Haskell at HOPL III

More recently...

- Paul Hudak, John Hughes, Simon Peyton Jones, Philip Wadler. *A History of Haskell: Being Lazy with Class*. In Proceedings of the Third ACM SIGPLAN 2007 Conference on History of Programming Languages (HOPL III), (San Diego, California, June 09 - 10, 2007), 12-1 - 12-55.

Check out the ACM Digital Library (www.acm.org/dl) for this article!

Concluding Course Meeting

- Thu, May 27, 2010: No lecture
- Thu, June 3, 2010: No lecture (public holiday)
- Thu, June 10, 2010: No lecture (epilog)
- Thu, June 10, 2010: possibly in the morning (depending on the availability of lecture rooms). More information will be posted asap on the homepage.

Last but not least

Final (oral) examination...

- In principle, any time. Just make an appointment by email (knoop@complang.tuwien.ac.at) or phone (58801-18510).
- Topics: Assignments and lecture materials.