

Implementation of Languages

SIMPLESEM

specify semantics in an abstract assembler language

static structure: data memory D, code memory C, instruction pointer ip

instructions:

- set target, source
- set 10, D[20]
- set 15, read
- set write, D[50]
- set 99, $D[15] + D[33 * C[41]] - 8$
- set D[10], D[20]
- jump 47
- jump 47, $D[3] > D[8]$
- jump D[13]
- halt

Language with Simple Statements

```

main()
{
    int i, j;           code memory           data memory
                        +-----+           +-----+
    get(i, j);          |0  set 0, read        |0  cell for i
                        |1  set 1, read        |1  cell for j
    while (i != j)      |2  jumpt 8, D[0] = D[1]
        if (i > j)      |3  jumpt 6, D[0] <= D[1]
            i -= j;     |4  set 0, D[0] - D[1]
        else           |5  jump 7
            j -= i;     |6  set 1, D[1] - D[0]
                        |7  jump 2
    print(i);           |8  set write, D[0]
                        |9  halt
}

```

Simple Routines

```
int i=1, j=2, k=3;
```

```
alpha()
```

```
{
```

```
    int i=4, l=5;
```

```
    ...
```

```
    i += k + l;
```

```
    ...
```

```
};
```

```
beta()
```

```
{
```

```
    int k=6;
```

```
    ...
```

```
    i = j + k;
```

```
    alpha();
```

```
    ...
```

```
};
```

```
main()
```

```
{
```

```
    ...
```

```
    beta();
```

```
    ...
```

```
}
```

Memory Layout (Simple Routines)

code memory

```

+-----+
|          // main
|014  set 6, 16
|015  jump 100
|          ...
|049  halt
|          // alpha
|058  set 4, D[4]+D[2]+D[5]
|          ...
|099  jump D[3]
|          // beta
|122  set 0, D[1]+D[7]
|123  set 3, 125
|124  jump 50
|          ...
|149  jump D[6]
|          // ...

```

data memory

```

+-----+
|0  8  //i    global
|1  2  //j
|2  3  //k
|3 125 //ret. alpha
|4 12  //i
|5  5  //l
|6 16  //ret. beta
|7  6  //k
|  ...

```

ip = 59

Separate Compilation

```
/* file 1 */
```

```
int i=1, j=2, k=3;  
extern beta();  
main()  
{  
    ...  
    beta();  
    ...  
}
```

```
/* file 2 */
```

```
extern int k;  
alpha()  
{  
    int i=4, l=5;  
    ...  
    i += k + l;  
    ...  
}
```

```
/* file 3 */
```

```
extern int i,j;  
extern alpha();  
beta()  
{  
    int k = 6;  
    ...  
    i = j + k;  
    alpha();  
    ...  
}
```

Recursion and Result Values

```
int n;
int fact()
{
    int loc;
    if (n > 1) {
        loc = n--;
        return loc * fact();
    }
    else
        return 1;
}
```

```
main()
{
    get(n);
    if (n >= 0)
        print(fact());
    else
        print("input error");
}
```

Stack with Dynamic Link

data memory

```
+-----  
|0  B    //CURRENT  
|1  C    //FREE  
|    ...           stack grows downwards  
|                   holds activation records  
|  
|A  ...  //return pointer (caller's activation record)  
|    ...  //dynamic link  
|    ...  //local variables  
|  
|B  ...  //return pointer (current activation record)  
|    A    //dynamic link  
|    ...  //local variables  
|  
|C  ...  //free memory  
|
```


Invocation and Return

set 1, D[1]+1	allocate memory for result
set D[1], ip+4	set return pointer
set D[1]+1, D[0]	set dynamic link
set 0, D[1]	set CURRENT
set 1, D[1]+AR	set FREE
jump addr	jump to routine
set 1, D[0]	set FREE
set 0, D[D[0]+1]	set CURRENT
jump D[D[1]]	jump to stored return pointer

Nested Blocks

int f() {	activation record of f
int x,y,w; //1	+-----
while(...) {	return pointer
int x,z; //2	dynamic link
while(...) {	x in //1
int y; ... //3	y in //1
}	w in //1
if(...) {	x in //2, a in //5
int x,w; ... //4	z in //2, b in //5
}	y in //3, x in //4, c in //5
}	w in //4, d in //5
if(...) {	+-----
int a,b,c,d; //5	
...	
}	
}	

Nested Routines and Static Link

```

int x,y,z;
void f1() {
    int t,u;
    void f2() {
        int x,w;
        void f3() {
            int y,w,t;
            ...
        }
        ...
    }
    ...
}

```

```

sketch of run-time stack
+-----
|A  global environment (x,y,z),
|    no dyn.link, stat.link needed
|B  act.record for f1 (t,u) with
|    dyn.link: A, stat.link: A
|C  act.record for f2 (x,w) with
|    dyn.link: B, stat.link: B
|D  act.record for f3 (y,w,t) with
|    dyn.link: C, stat.link: C
|E  act.record for f2 (x,w) with
|    dyn.link: D, stat.link: B !!
|

```

Computing Frame Pointer

each variable statically bound to a pair $\langle d, o \rangle$ where

d is the **distance** between current and addressed static scope

e.g., $d=2$ when accessing u (defined in f_1) within f_3 ,

o is the **offset** of the variable within the activation record

computing the **frame pointer**:

$$fp(d) = \text{if } d=0 \text{ then } D[0] \text{ else } D[fp(d-1)+stat.link.offset]$$

examples: $fp(0) = D[0]$,

$fp(1) = D[D[0]+stat.link.offset]$,

$fp(2) = D[D[D[0]+stat.link.offset]+stat.link.offset]$

adresse of variable bound to $\langle d, o \rangle$ is $D[fp(d)+o]$

Routine Invocation with Static Link

set 1, D[1]+1	allocate memory for result
set D[1], ip+5	set return pointer
set D[1]+1, D[0]	set dynamic link
set D[1]+2, fp(d)	set static link (stat.link.offset = 2)
set 0, D[1]	set CURRENT
set 1, D[1]+AR	set FREE
jump addr	jump to routine

Dynamic Array

example in Ada: `type V is array (Integer range <>) of Integer;`
`A: V (1 .. N);` -- N not statically known

compiler allocates in activation record space for **descriptor**
(holds index range and pointer to memory of array)

when executing the array declaration at run-time:

extend current activation record with space for array,
initialize descriptor appropriately

example for access `A(I) := 0;`

`set[D[D[0]+m] + D[D[0]+s]], 0`

(m = offset of descriptor of A; s = offset of local variable I)

Dynamic Scoping

void sub2() {	sketch of run-time stack
declare x;	+-----
...	A link = none (act.record for main)
}	x = ...
void sub1() {	y = ...
declare y;	z = ...
sub2();	B link = A (act.record for sub1)
...	y = ...
}	C link = B (act.record for sub2)
void main() {	z = ...
declare x,y,z;	
sub1();	
sub2();	dynamic search for name along link
}	

Heap

local data on stack are lost when returning from a routine invocation

non-local data needed after return must be allocated on the **heap**

data memory

```
+-----  
|  stack: grows downwards,  
|           holds activation records removed on return  
|  
|  
|  heap:  grows upwards,  
|           holds data alive until program termination  
+-----
```


Parameter Passing

call by reference: pass l-value of actual parameter to callee

call by copy: parameter is local variable of callee

call by value: copy r-value into variable of callee

call by result: copy result back into variable of caller

call by value-result:] call by value on invocation
and call by result on return (value copied twice)

call by name: substitute actual parameter
for each occurrence of formal parameter in callee

Implementation of Call by Reference

formal parameter is reference to actual parameter in callee's activation record
(offset off)

pass l-value if actual parameter is a variable $\langle d, o \rangle$ in caller:

set $D[0] + \text{off}$, $\text{fp}(d) + o$

pass r-value if actual parameter is call-by-reference parameter $\langle d, o \rangle$ in caller:

set $D[0] + \text{off}$, $D[\text{fp}(d) + o]$

assignment to a call-by-reference parameter: $x = 0$;

set $D[D[0] + \text{off}]$, 0

Call by Reference vs. Value Result

these calling conventions differ if

two formal parameters are aliases

example: actual parameters: $a[i]$ und $a[j]$
formal parameters: x und y
body of routine: $x = 0; y++;$

a formal parameter and a variable visible to both (caller and callee) are alias

example: actual parameter: a
formal parameter: x
body of routine: $a = 1; x = x + a;$

Call by Name

```
swap(int a, b) {  
    int temp = a;  
    a = b;  
    b = temp;  
};  
...  
swap(i, a[i]);  
...  
  
temp = i;  
i = a[i];  
a[i] = temp;  
  
|   int c; /* global variable */  
|   swap(int a, b) {  
|       int temp = a; a = b;  
|       b = temp; c++;  
|   };  
|   y() {  
|       int c, d;  
|       swap(c, d);  
|   };  
  
|   temp = c; c = d;  
|   d = temp;  
|   c++;           !!
```

Routine as Parameter

```

1  int u, v;      |    7  void b(routine x) |   19  void main()
2  void a()       |    8  {                  |   20  {
3  {              |    9      int u, v, y;    |   21      b(a);
4      int y;      |   10      void c() {      |   22  };
5      ...         |   11          ...         |
6  };             |   12          y = ...;    |
                  |   13          ...         |
                  |   14      };             |
                  |   15      x();           |
                  |   16      b(c);          |
                  |   17      ...           |
                  |   18  }                 |

```

routine parameter must include static link